

I631: Foundation of Computational Geometry (13) Envelopes and Levels I

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Goal of this lecture

Background

- A hyperplane arrangement has abundant information
- Often, we're interested in substructures of a hyperplane arrangement
- Envelopes and levels are examples of such structures

Goal of this lecture

- Learn the relevant notions for envelopes and levels
- Learn the connection with Voronoi diagrams

- ① Voronoi diagrams
- ② Envelopes and levels
- ③ Relationship with Voronoi diagrams

Nearest neighbors

$P = \{p_1, \dots, p_n\} \subseteq \mathbb{R}^d$ a finite point set

Def.: Nearest neighbor

A **nearest neighbor** of a point $q \in \mathbb{R}^d$ in P is a point $p \in P$ such that

$$d(p, q) \leq d(p', q) \quad \forall p' \in P$$

Voronoi diagrams

$P = \{p_1, \dots, p_n\} \subseteq \mathbb{R}^d$ a finite point set, $n \geq 2$

Def.: Voronoi diagram

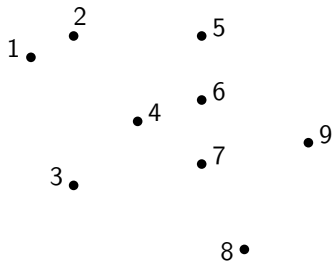
The **Voronoi diagram** of P is a partition of \mathbb{R}^d by the regions

$$\text{vor}(S) = \{q \in \mathbb{R}^d \mid S = \text{the nearest neighbors of } q \text{ in } P\}$$

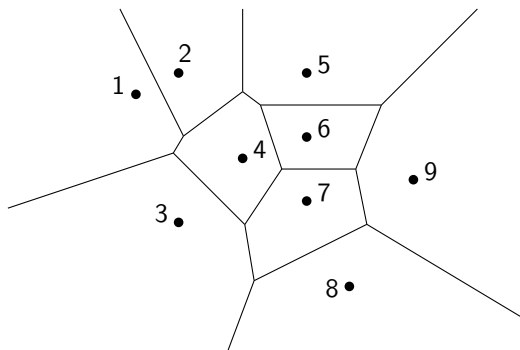
for all $S \subseteq P$, $|S| \geq 1$

- Denote the Voronoi diagram of P by $\text{Vor}(P)$
- Each non-empty $\text{vor}(S)$ is called the **Voronoi region** of S

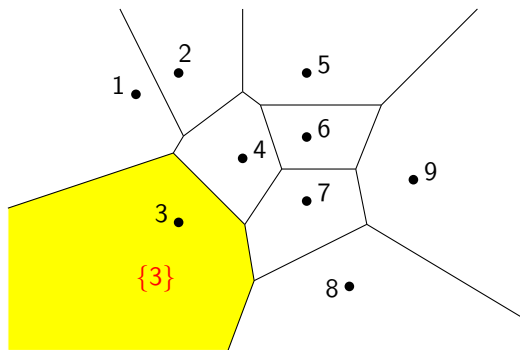
Example: Voronoi diagrams



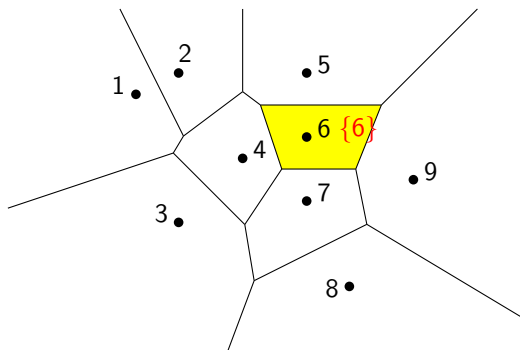
Example: Voronoi diagrams



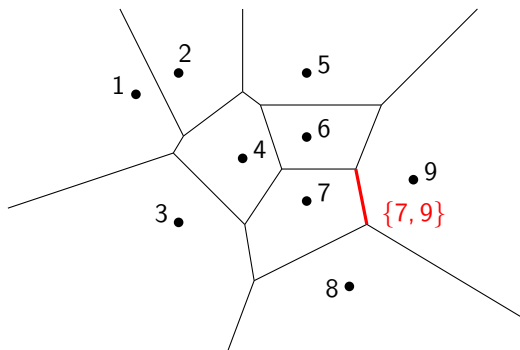
Example: Voronoi diagrams



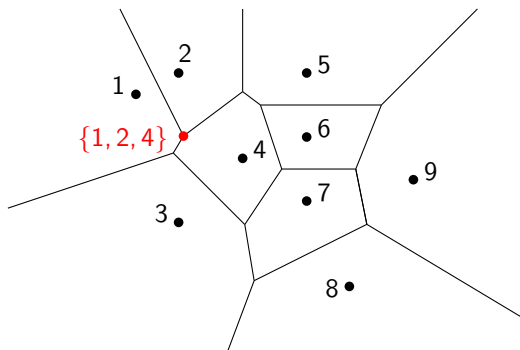
Example: Voronoi diagrams



Example: Voronoi diagrams



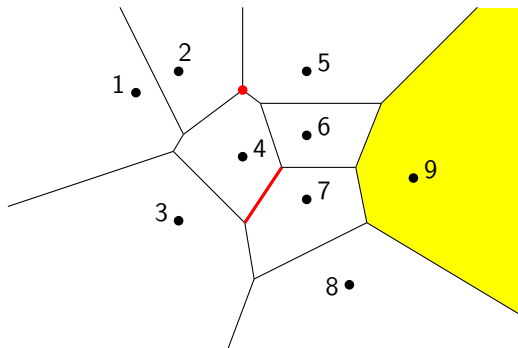
Example: Voronoi diagrams



Dimension of Voronoi regions

Def.: Dimension of a Voronoi region

The **dimension** of $\text{vor}(S)$ is
the dimension of a minimal affine subspace containing $\text{vor}(S)$



Special Voronoi regions

Voronoi regions have names according to their dimensions

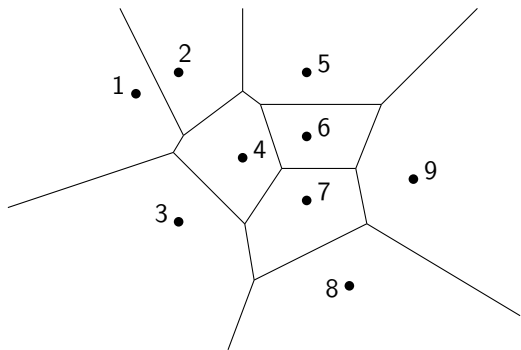
- Voronoi vertex: 0-dimensional Voronoi region
- Voronoi edge: 1-dimensional Voronoi region
- Voronoi ridge: $d-2$ -dimensional Voronoi region
- Voronoi facet: $d-1$ -dimensional Voronoi region
- Voronoi cell: d -dimensional Voronoi region

The number of Voronoi cells

Question

How many Voronoi cells can there be in the Voronoi diagram of a set of n points in \mathbb{R}^d ?

This determines the intrinsic difficulty of the problem of computing the Voronoi diagram of a given point set



2-Nearest neighbors

$P = \{p_1, \dots, p_n\} \subseteq \mathbb{R}^d$ a finite point set

Def.: 2-Nearest neighbor

A **2-nearest neighbor** of a point $q \in \mathbb{R}^d$ in P is a point $p \in P$ such that there exists a set X with $|X| \leq 1$ and

$$d(p, q) \leq d(p', q) \quad \forall p' \in P \setminus X$$

Order-2 Voronoi diagrams

$P = \{p_1, \dots, p_n\} \subseteq \mathbb{R}^d$ a finite point set, $n \geq 3$

Def.: Order-2 Voronoi diagram

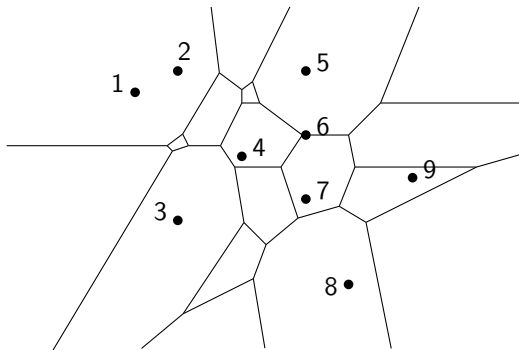
The **order-2 Voronoi diagram** of P is a partition of \mathbb{R}^d by the regions

$$\text{vor}^{(2)}(S) = \{q \in \mathbb{R}^d \mid S = \text{the 2-nearest neighbors of } q \text{ in } P\}$$

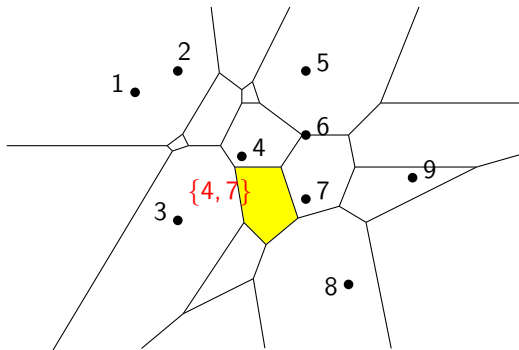
for all $S \subseteq P$, $|S| \geq 2$

- Denote the order-2 Voronoi diagram of P by $\text{Vor}^{(2)}(P)$
- Each non-empty $\text{vor}^{(2)}(S)$ is called the **Voronoi region** of S
 - Dimensions and special names are defined similarly

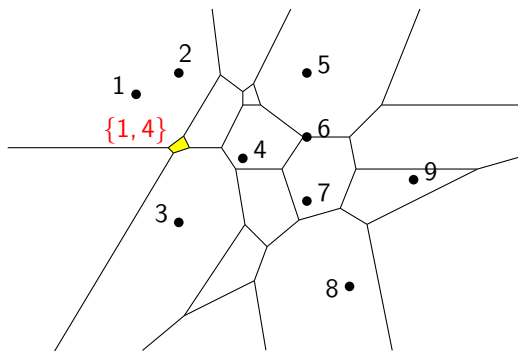
Example: Order-2 Voronoi diagrams



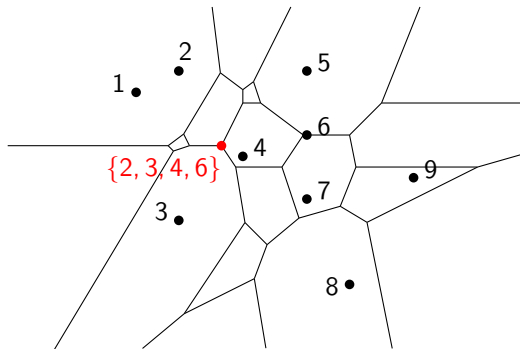
Example: Order-2 Voronoi diagrams



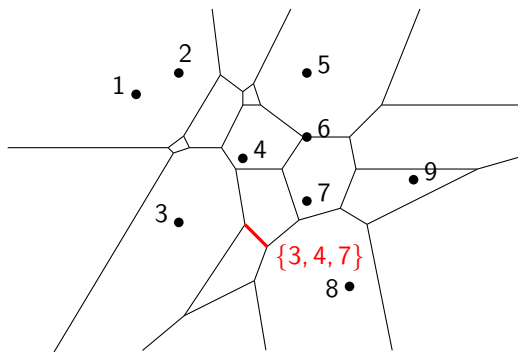
Example: Order-2 Voronoi diagrams



Example: Order-2 Voronoi diagrams



Example: Order-2 Voronoi diagrams



k -Nearest neighbors

$P = \{p_1, \dots, p_n\} \subseteq \mathbb{R}^d$ a finite point set

Def.: k -Nearest neighbor

A **k -nearest neighbor** of a point $q \in \mathbb{R}^d$ in P is a point $p \in P$ such that there exist a set $X \subseteq P$ with $|X| \leq k-1$ and

$$d(p, q) \leq d(p', q) \quad \forall p' \in P \setminus X$$

Order- k Voronoi diagrams

$P = \{p_1, \dots, p_n\} \subseteq \mathbb{R}^d$ a finite point set, $n \geq k+1$

Def.: Order- k Voronoi diagram

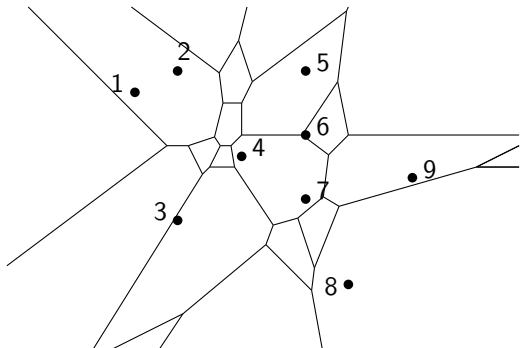
The **order- k Voronoi diagram** of P is a partition of \mathbb{R}^d by the regions

$$\text{vor}^{(k)}(S) = \{q \in \mathbb{R}^d \mid S = \text{the } k\text{-nearest neighbors of } q \text{ in } P\}$$

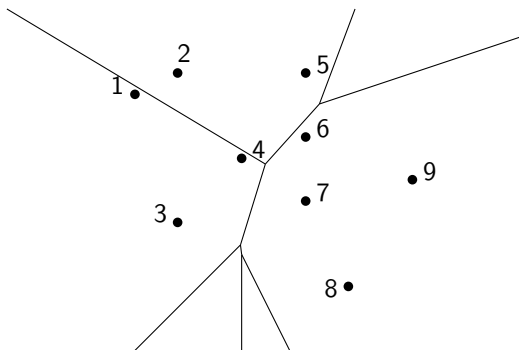
for all $S \subseteq P$, $|S| \geq k$

- Denote the Voronoi diagram of P by $\text{Vor}^{(k)}(P)$
- Each non-empty $\text{vor}^{(k)}(S)$ is called the **Voronoi region** of S
 - Dimensions and special names are defined similarly

Example: Order-3 Voronoi diagrams



Example: Farthest-point Voronoi diagrams



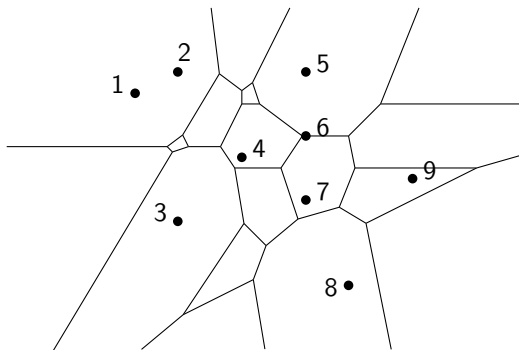
The order- $(n-1)$ Voronoi diagram is usually called the **farthest-point Voronoi diagram**

The number of Voronoi cells in $\text{Vor}^{(k)}(P)$

Question

How many can Voronoi cells there be in the order- k Voronoi diagram of a set of n points in \mathbb{R}^d ?

This determines the intrinsic difficulty of the problem of computing the order- k Voronoi diagram of a given point set



The rest of today's lecture

We will see

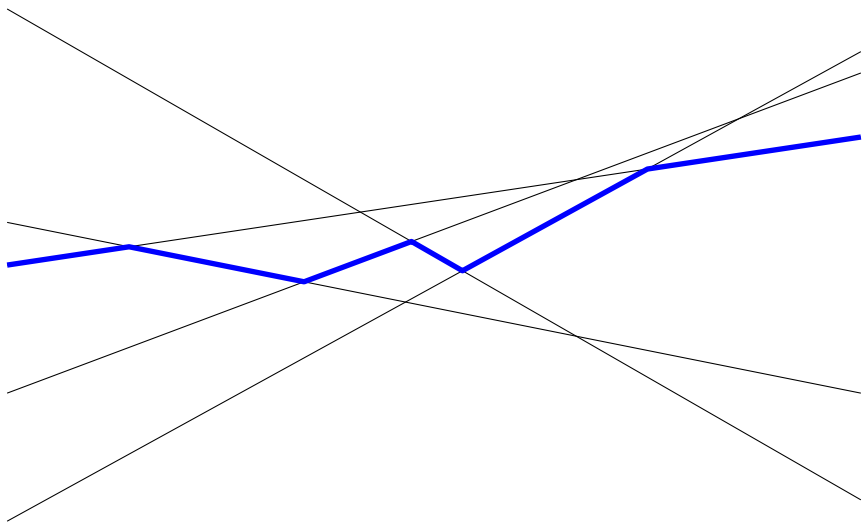
Voronoi diagrams and order- k Voronoi diagrams are closely related to envelopes and levels of hyperplane arrangements

The rest of the lecture

- Definition of envelopes and levels
- Relationship with Voronoi diagrams

- ① Voronoi diagrams
- ② Envelopes and levels
- ③ Relationship with Voronoi diagrams

Example: Levels of a hyperplane arrangement



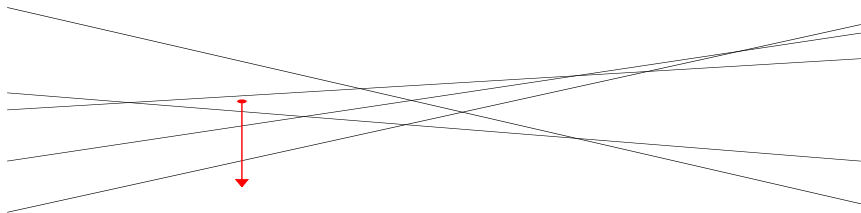
Levels of a point in a hyperplane arrangement

$\mathcal{A} = \{H_1, \dots, H_n\}$ a hyperplane arrangement in \mathbb{R}^d where
 $H_i = \{x \in \mathbb{R}^d \mid a_i \cdot x = b_i\}$, $a_i \in \mathbb{R}^d \setminus \{0\}$ ($a_i \cdot e_d > 0$), $b_i \in \mathbb{R}$

Level of a point

The **level** of a point $p \in \mathbb{R}^d$ in \mathcal{A} is
 the number of hyperplanes in \mathcal{A} below p ;
 Alternatively, the level of p is k if

$$k = |\{i \in \{1, \dots, n\} \mid a_i \cdot p < b_i\}|$$

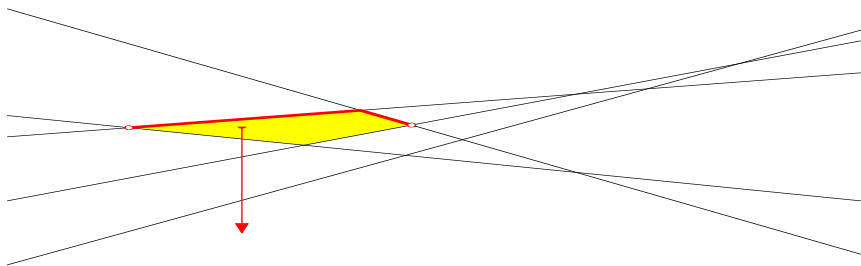


Levels of a point in a hyperplane arrangement

$\mathcal{A} = \{H_1, \dots, H_n\}$ a hyperplane arrangement in \mathbb{R}^d where
 $H_i = \{x \in \mathbb{R}^d \mid a_i \cdot x = b_i\}$, $a_i \in \mathbb{R}^d \setminus \{0\}$ ($a_i \cdot e_d > 0$), $b_i \in \mathbb{R}$

Observation

Every point in the same face in the arrangement has the same level

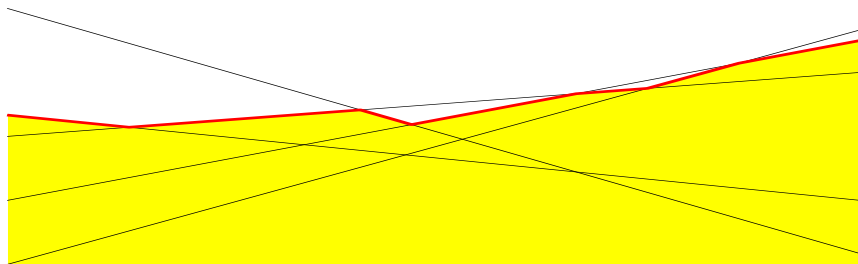


Levels of a hyperplane arrangement

$\mathcal{A} = \{H_1, \dots, H_n\}$ a hyperplane arrangement in \mathbb{R}^d where
 $H_i = \{x \in \mathbb{R}^d \mid a_i \cdot x = b_i\}$, $a_i \in \mathbb{R}^d \setminus \{0\}$ ($a_i \cdot e_d > 0$), $b_i \in \mathbb{R}$

Def.: Levels of a hyperplane arrangement

The **k -level** of \mathcal{A} is the boundary of the set of all points of level at most k

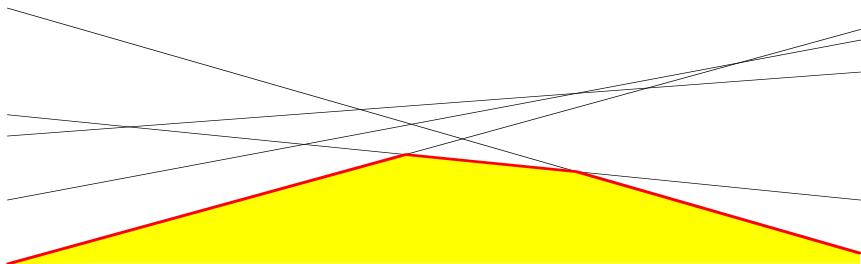


Lower envelope of a hyperplane arrangement

$\mathcal{A} = \{H_1, \dots, H_n\}$ a hyperplane arrangement in \mathbb{R}^d where
 $H_i = \{x \in \mathbb{R}^d \mid a_i \cdot x = b_i\}$, $a_i \in \mathbb{R}^d \setminus \{0\}$ ($a_i \cdot e_d > 0$), $b_i \in \mathbb{R}$

Def.: Lower envelope of a hyperplane arrangement

The **lower envelope** of \mathcal{A} is the 0-level of \mathcal{A}

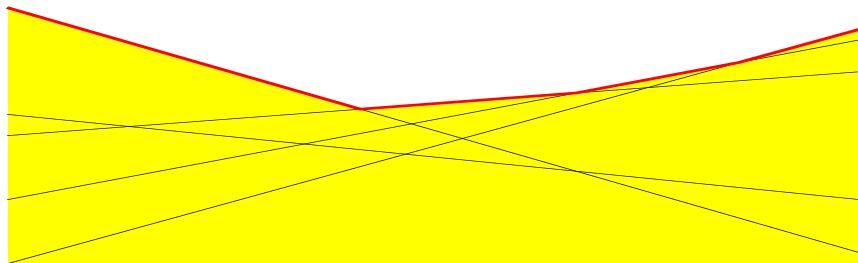


Upper envelope of a hyperplane arrangement

$\mathcal{A} = \{H_1, \dots, H_n\}$ a hyperplane arrangement in \mathbb{R}^d where
 $H_i = \{x \in \mathbb{R}^d \mid a_i \cdot x = b_i\}$, $a_i \in \mathbb{R}^d \setminus \{0\}$ ($a_i \cdot e_d > 0$), $b_i \in \mathbb{R}$

Def.: Upper envelope of a hyperplane arrangement

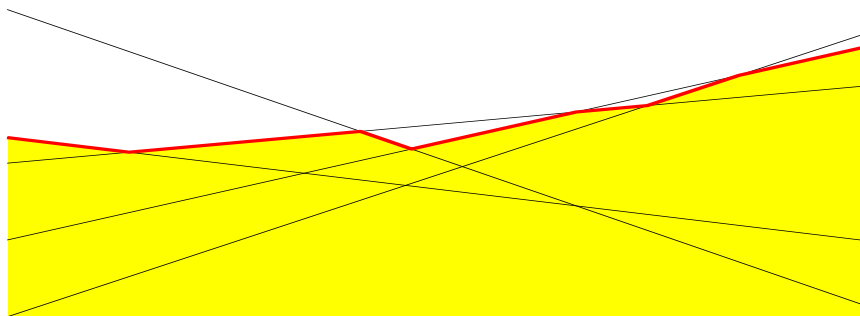
The **upper envelope** of \mathcal{A} is the $(n-1)$ -level of \mathcal{A}



Structures of levels

The k -level of a hyperplane arrangement is a collection of polyhedra

- They are faces of the hyperplane arrangements
- In particular, they are convex

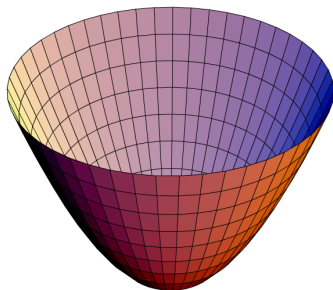
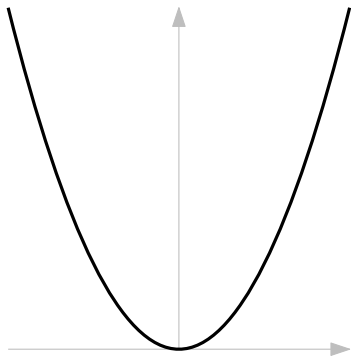


- ① Voronoi diagrams
- ② Envelopes and levels
- ③ Relationship with Voronoi diagrams**

Unit paraboloids

Define the **unit paraboloid** in \mathbb{R}^{d+1} as

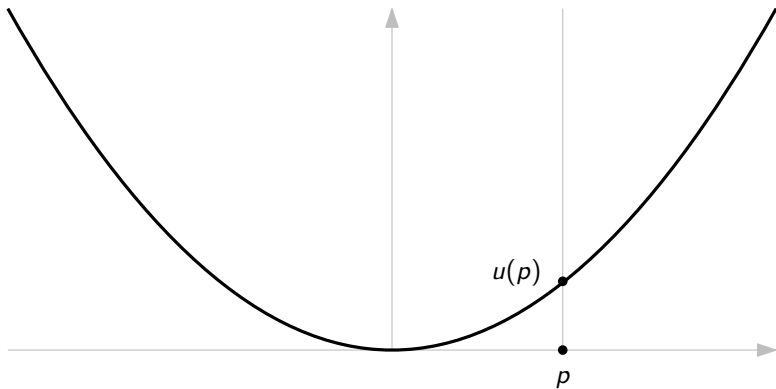
$$U = \{x \in \mathbb{R}^{d+1} \mid x_{d+1} = x_1^2 + x_2^2 + \cdots + x_d^2\}$$



<http://en.wikipedia.org/wiki/Paraboloid>

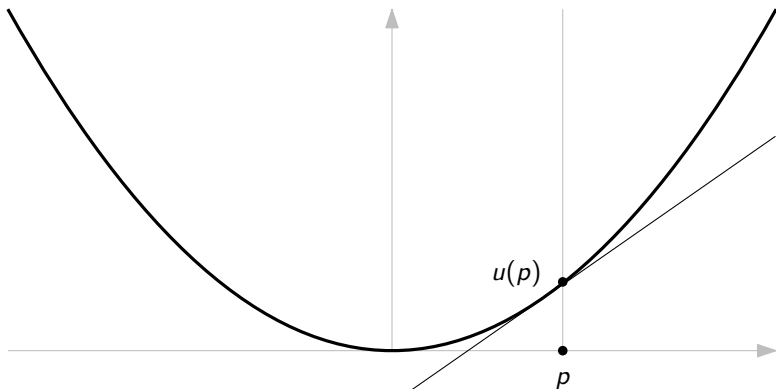
Lifting up a point set to the unit paraboloid

For $p \in \mathbb{R}^d$, let $u(p) = (p, p_1^2 + p_2^2 + \cdots + p_d^2) \in \mathbb{R}^{d+1}$



Tangent hyperplanes

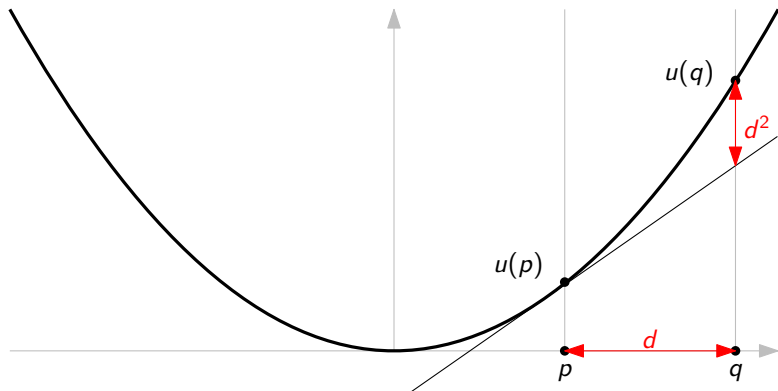
The hyperplane $H(p) = \{x \in \mathbb{R}^{d+1} \mid x_{d+1} = \sum_{i=1}^d (2p_i x_i - p_i^2)\}$ is tangent to U at $u(p)$ (Exercise)



Relationship with distances

Observation

For points $p, q \in \mathbb{R}^d$,
the vertical distance from q to $H(p)$ is $d(p, q)^2$



Proof of the observation

- $d(p, q)^2 = \sum_{i=1}^d (p_i - q_i)^2$

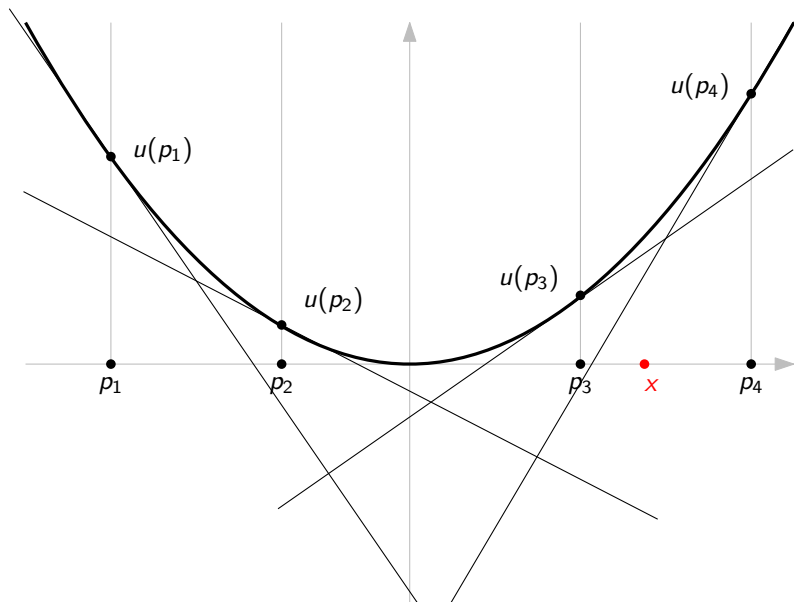
- The vertical distance from q to $H(p)$ is

$$\begin{aligned}
 &= u(q) - \sum_{i=1}^d (2p_i q_i - p_i^2) \\
 &= \sum_{i=1}^d q_i^2 - \sum_{i=1}^d (2p_i q_i - p_i^2) \\
 &= \sum_{i=1}^d (q_i^2 - 2p_i q_i + p_i^2) = \sum_{i=1}^d (p_i - q_i)^2
 \end{aligned}$$

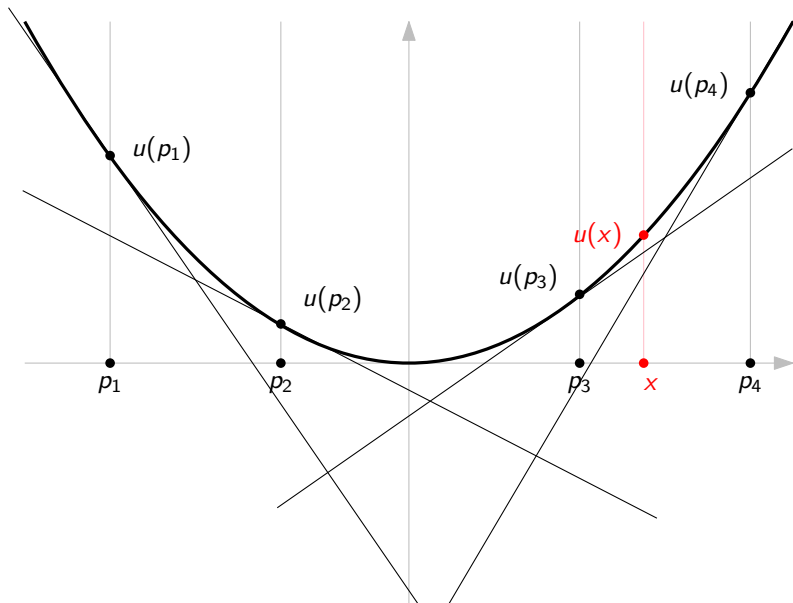
- These two are equal



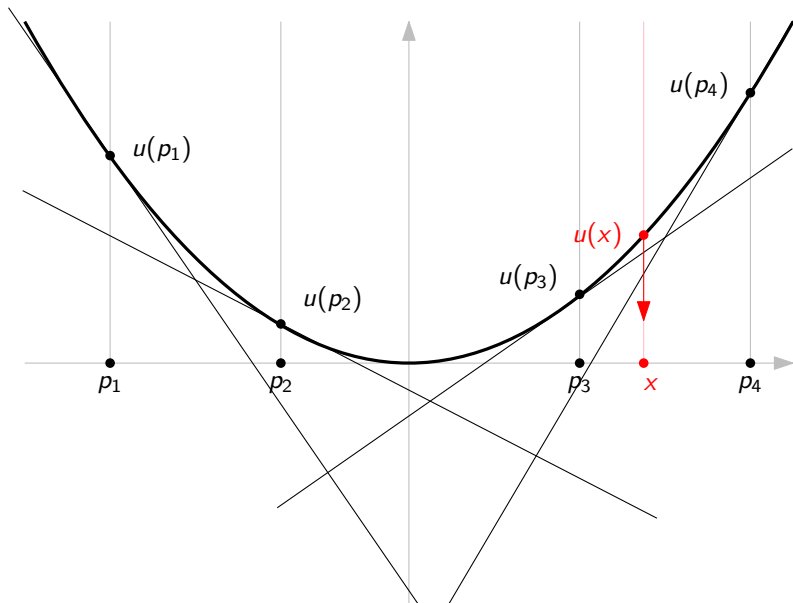
Comparing distances



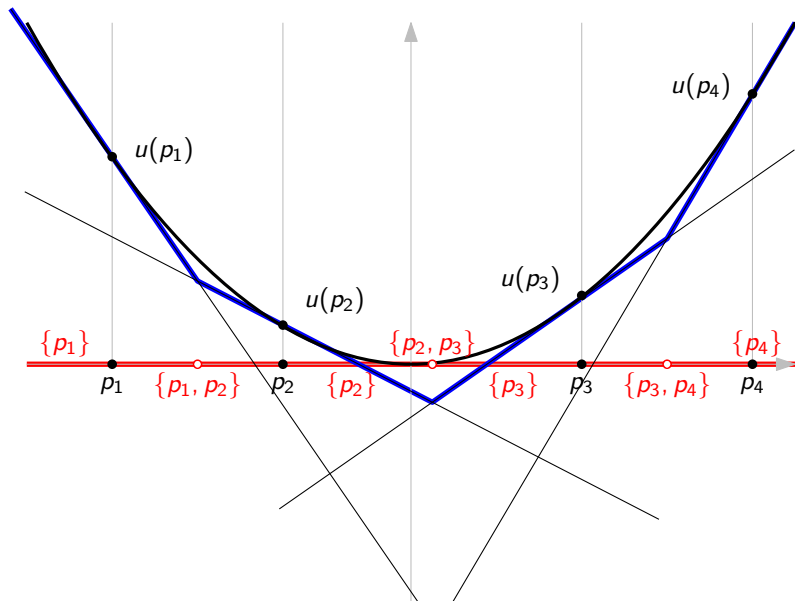
Comparing distances



Comparing distances



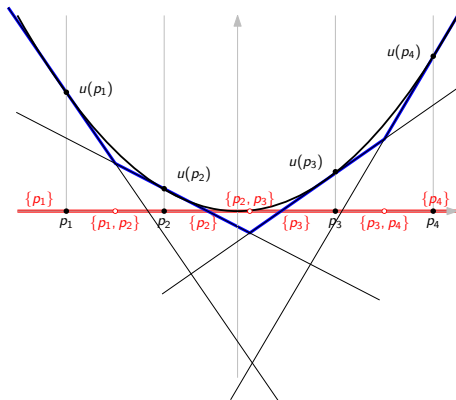
Comparing distances



Relationship with Voronoi diagrams

Proposition

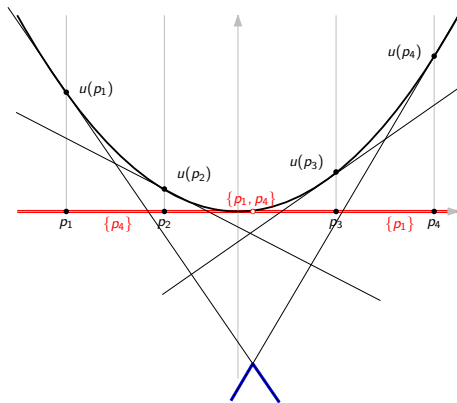
The Voronoi diagram of P is the projection of the upper envelope of the hyperplane arrangement $\{H(p) \mid p \in P\}$ to the hyperplane $\{x \in \mathbb{R}^{d+1} \mid x_{d+1} = 0\}$



Relationship with farthest-point Voronoi diagrams

Proposition

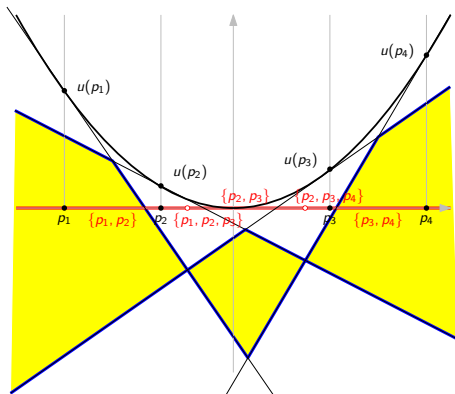
The farthest-point Voronoi diagram of P is the projection of the lower envelope of the hyperplane arrangement $\{H(p) \mid p \in P\}$ to the hyperplane $\{x \in \mathbb{R}^{d+1} \mid x_{d+1} = 0\}$



Relationship with order-2 Voronoi diagrams

Fact

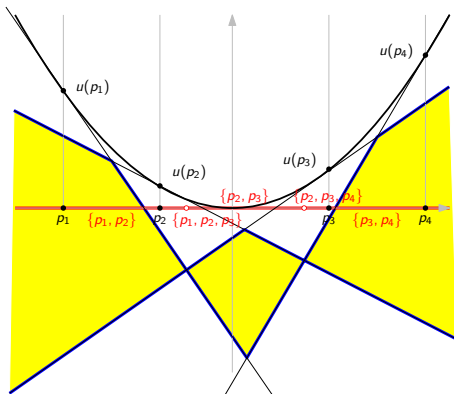
The order-2 Voronoi diagram of P is the projection of the part of the hyperplane arrangement $\{H(p) \mid p \in P\}$ between the $(n-2)$ -level and the $(n-3)$ -level to the hyperplane $\{x \in \mathbb{R}^{d+1} \mid x_{d+1} = 0\}$



Relationship with order- k Voronoi diagrams

Fact

The order- k Voronoi diagram of P is the projection of the part of the hyperplane arrangement $\{H(p) \mid p \in P\}$ between the $(n-k)$ -level and the $(n-k-1)$ -level to the hyperplane $\{x \in \mathbb{R}^{d+1} \mid x_{d+1} = 0\}$



The number of Voronoi cells in Voronoi diagrams

Consequence

The Voronoi diagram of a set of n points in \mathbb{R}^d has

- at most n Voronoi cells
- at most $O(n^{\lceil d/2 \rceil})$ vertices

This is a consequence of the facts that

- The upper envelope is the boundary of a polyhedron with at most n facets
- The number of vertices of a d -dim polytope with n facets is $O(n^{\lfloor d/2 \rfloor})$ (Upper bound theorem)

The number of Voronoi cells in farthest-point Voronoi diagrams

Consequence

The farthest-point Voronoi diagram of a set of n points in \mathbb{R}^d has

- at most n Voronoi cells
- at most $O(n^{\lceil d/2 \rceil})$ vertices

This is a consequence of the facts that

- The lower envelope is the boundary of a polyhedron with at most n facets
- The number of vertices of a d -dim polytope with n facets is $O(n^{\lfloor d/2 \rfloor})$ (Upper bound theorem)

Convexity of Voronoi regions

Consequence

Every Voronoi region of the order- k Voronoi diagram of a point set in \mathbb{R}^d is convex

This is a consequence of the facts that

- Each face of the hyperplane arrangement is convex
- The projection of a convex set is convex (Exercise)

Voronoi diagrams

- Def.: A partition of \mathbb{R}^d with respect to the distances to a given point set
- Terms: Voronoi regions, Voronoi cells, ...

Levels and Envelopes

- Level of a point: the number of hyperplanes below the point
- k -Level: the boundary of the set of points with level at most k
- Lower envelope: 0-level
- Upper envelope: $(n-1)$ -level

Connection to Voronoi diagrams: Lift up to the unit paraboloid, and consider the hyperplane arrangement induced by tangents

Fact (Lee '82)

The order- k Voronoi diagram of a set of n points in \mathbb{R}^2 has $O(k(n - k))$ Voronoi cells; This bound is tight

Not much is known for the number of Voronoi cells in the order- k Voronoi diagrams when $d \geq 3$

- Even in \mathbb{R}^2 , determining the maximum number of edges in the k -level of a line arrangement is a difficult problem
- We'll look at an argument that gives some upper bound

- Matoušek: *Lectures on Discrete Geometry*
 - 4, 5, 11
- Edelsbrunner: *Algorithms in Combinatorial Geometry*
 - Chapters 1, 13

- ① Voronoi diagrams
- ② Envelopes and levels
- ③ Relationship with Voronoi diagrams