

## Today's contents

# Topics on Computing and Mathematical Sciences I Graph Theory (11) Extremal Graph Theory II

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① Review of Lecture 9

② Density and  $\varepsilon$ -regular pairs

③ Szemerédi's regularity lemma

④ Embedding lemma and a proof of the Erdős-Stone theorem

⑤ Algorithmic issue

A complete answer for the Turán-type problems

Theorem 9.3 (Erdős, Simonovits '66; Erdős-Simonovits-Stone thm)

$\forall$  graph  $H$

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = 1 - \frac{1}{\chi(H) - 1};$$

In other words,

$$\text{ex}(n, H) = \left(1 - \frac{1}{\chi(H) - 1}\right) \binom{n}{2} + o(n^2)$$

Proof outline I (Lecture 9)

① Deduce the Erdős-Simonovits-Stone thm from the Erdős-Stone thm

② Deduce the Erdős-Stone thm from its weaker version

③ Prove the weaker version of Erdős-Stone thm

Proof outline II (today)

① Deduce the Erdős-Simonovits-Stone thm from the Erdős-Stone thm

② Use Szemerédi's regularity lemma to prove the embedding lemma

③ Use the embedding lemma to prove Erdős-Stone thm

## Erdős-Stone theorem

Theorem 9.4 (Erdős, Stone '46)

 $\forall s \geq 1, r \geq 2$  natural numbers

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, T_{rs,r})}{\binom{n}{2}} = 1 - \frac{1}{r-1}$$

We have seen the Erdős-Stone theorem 9.4 implies the Erdős-Stone-Simonovits theorem 9.3

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④ Embedding lemma and a proof of the Erdős-Stone theorem

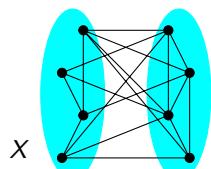
⑤ Algorithmic issue

## Density

## Notation

 $G = (V, E)$  a graph;  $X, Y \subseteq V$  disjoint

- $e(X, Y) = |\{ \{u, v\} \in E \mid u \in X, v \in Y \}|$
- $d(X, Y) = \frac{e(X, Y)}{|X||Y|}$  (called the **density** of  $(X, Y)$ )



$$e(X, Y) = 10$$

$$d(X, Y) = \frac{10}{4 \times 4} = \frac{5}{8}$$

 $\varepsilon$ -Regular pairs $G = (V, E)$  a graph;  $A, B \subseteq V$  disjoint;  $\varepsilon > 0$ Definition ( $\varepsilon$ -Regular pair) $(A, B)$  is  $\varepsilon$ -regular if for all  $X \subseteq A$  and  $Y \subseteq B$ ,

$$|X| \geq \varepsilon |A|, |Y| \geq \varepsilon |B| \Rightarrow |d(X, Y) - d(A, B)| \leq \varepsilon$$

## Intuition

Large subsets  $X, Y$  of  $A, B$  (resp.) look similar to  $A, B$  (as far as the density is concerned)

## Regular pairs look like random bipartite graphs (1)

## Lemma 11.1

$G = (V, E)$  bipartite w/ its partite sets  $A, B$ ,  $|A| = |B| = n$ ;  
 $(A, B)$   $\varepsilon$ -regular w/ its density  $d \Rightarrow$

$$|\{v \in A \mid (d-\varepsilon)n \leq d(v) \leq (d+\varepsilon)n\}| \geq (1-2\varepsilon)n.$$

## Note

Construct a bipartite random graph  $G = (A \cup B, E)$  by joining two vertices  $v \in A, u \in B$  w/ prob.  $d/n^2$  independently at random. Then

- $\mathbb{E}[d(A, B)] = d$
- $\mathbb{E}[d(v)] = dn$

## Proof idea of Lemma 11.1

Suppose not

- $A_1 = \{v \in A \mid d(v) < (d-\varepsilon)n\}$ 
  - $\therefore d(A_1, B) < d-\varepsilon$
- $A_2 = \{v \in A \mid d(v) > (d+\varepsilon)n\}$  (Note:  $A_1 \cap A_2 = \emptyset$ )
  - $\therefore d(A_2, B) > d+\varepsilon$
- Then,  $|A_1 \cup A_2| \geq 2\varepsilon n$
- $\therefore |A_1| \geq \varepsilon n$  or  $|A_2| \geq \varepsilon n$
- $|A_1| \geq \varepsilon n \Rightarrow |d(A_1, B) - d(A, B)| \leq \varepsilon$  (by  $\varepsilon$ -regularity of  $(A, B)$ )
- $|A_2| \geq \varepsilon n \Rightarrow |d(A_2, B) - d(A, B)| \leq \varepsilon$  (by  $\varepsilon$ -regularity of  $(A, B)$ )
- In both cases, we have a contradiction  $\square$

## Regular pairs look like random bipartite graphs (2)

## Lemma 11.2

## Setup

- $G = (V, E)$  a tripartite graph with its partite sets  $V_1, V_2, V_3$
- $|V_1| = |V_2| = |V_3| = n$
- $(V_i, V_j)$  is  $\varepsilon$ -regular with density  $d$ ,  $\forall 1 \leq i < j \leq 3$ , for some  $\varepsilon > 0$ ,  $0 < d < 1$
- $\varepsilon \leq d/2$

## Statement

- #  $K_3$ 's in  $G \geq (1-4\varepsilon)(d-\varepsilon)^3n^3$

Note: the exact coefficient " $(1-4\varepsilon)(d-\varepsilon)^3$ " is not important, but it's important to note the coefficient only depends on  $\varepsilon$  and  $d$

## Proof idea of the lemma 11.2 (1/2)

First, estimate the number of vertices in  $V_1$  that have many neighbors in both of  $V_2$  and  $V_3$

Let  $X_i = |\{v \in V_1 \mid (d-\varepsilon)n \leq |N_G(v) \cap V_i| \leq (d+\varepsilon)n\}|$  ( $i \in \{2, 3\}$ )

$$\begin{aligned} |X_2 \cap X_3| &= |X_2| + |X_3| - |X_2 \cup X_3| \\ &\geq (1-2\varepsilon)n + (1-2\varepsilon)n - n \\ &\quad (\text{by Lem 11.1}) \\ &= (1-4\varepsilon)n \end{aligned}$$

## Proof idea of Lemma 11.2 (2/2)

- Fix an arbitrary  $v \in X_2 \cap X_3$
- let  $S_i(v) = N_G(v) \cap V_i$  ( $i \in \{2, 3\}$ )
- $|S_2(v)| \geq (d-\varepsilon)n \geq (2\varepsilon-\varepsilon)n = \varepsilon n$
- $|S_3(v)| \geq (d-\varepsilon)n \geq (2\varepsilon-\varepsilon)n = \varepsilon n$  (similarly)
- $\therefore$  by the  $\varepsilon$ -regularity of  $(V_2, V_3)$ , we have

$$\begin{aligned} e(S_2(v), S_3(v)) &= d(S_2(v), S_3(v))|S_2(v)||S_3(v)| \\ &\geq (d-\varepsilon) \cdot (d-\varepsilon)n \cdot (d-\varepsilon)n \\ &= (d-\varepsilon)^3 n^2 \end{aligned}$$

- Hence

$$\begin{aligned} \# \text{ of } K_3 \text{'s in } G &\geq \sum_{v \in X_2 \cap X_3} e(S_2(v), S_3(v)) \\ &\geq (1-4\varepsilon)n \cdot (d-\varepsilon)^3 n^2 \\ &= (1-4\varepsilon)(d-\varepsilon)^3 n^3 \quad \square \end{aligned}$$

## Extension of the previous lemma

## Lemma 11.3

$\forall r \in \mathbb{N} \forall d \in (0, 1) \exists \varepsilon_0 < 1 \exists c > 0 \forall \varepsilon \leq \varepsilon_0 \forall \text{ graph } G \text{ s.t.}$

- $G$  an  $r$ -partite graph with its partite sets  $V_1, \dots, V_r$
- $|V_1| = \dots = |V_r| = n$
- $(V_i, V_j)$  is  $\varepsilon$ -regular with density  $d$ ,  $\forall i, j \in \{1, \dots, r\}$

it holds that

- $\# K_r$ 's in  $G \geq cn^r$

Proof: Exercise (Hint: Induction on  $r$ )

## Note

From the previous lemma, when  $r = 3$  we can take  $\varepsilon_0 = d/2$  and  $c = (1-4\varepsilon_0)(d-\varepsilon_0)^3$

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 $\varepsilon$ -Regular partitions

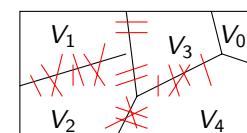
$G = (V, E)$  a graph;  $\varepsilon > 0$ ;

$\{V_0, V_1, \dots, V_k\}$  a partition of  $V$  (w/ possibly  $V_0 = \emptyset$ )

Definition ( $\varepsilon$ -Regular partition)

$\{V_0, V_1, \dots, V_k\}$  is  $\varepsilon$ -regular if

- $|V_0| \leq \varepsilon|V|$
- $|V_1| = \dots = |V_k|$
- $|\{(i, j) \mid 1 \leq i < j \leq k, (V_i, V_j) \text{ not } \varepsilon\text{-regular}\}| \leq \varepsilon k^2$



$V_0$  is called an **exceptional set**

## Szemerédi's regularity lemma

Szemerédi's regularity lemma (Szemerédi '76)

$\forall \varepsilon > 0$  a real number  $\forall n_0 \geq 1$  an integer

$\exists N_0 \in \mathbb{N} \forall G$  a graph:

$$n(G) \geq n_0 \implies \begin{aligned} G \text{ has an } \varepsilon\text{-regular partition} \\ \{V_0, V_1, \dots, V_k\} \text{ w/ } n_0 \leq k \leq N_0 \end{aligned}$$

### Intuition

Every dense graph looks like a union of random bipartite graphs

- Often called “Szemerédi’s uniformity lemma” too
- We are not going to prove Szemerédi’s regularity lemma, but looking at some applications of it

## Quantitative version of Mantel’s theorem

The following proposition gives an estimate for the number of  $K_3$ ’s contained in an  $n$ -vertex graph with more than  $n^2/4$  edges

**Proposition 11.4 (Quantitative version of Mantel’s theorem)**

$\forall \delta > 0 \exists n_0 \in \mathbb{N} \exists c > 0 \forall G = (V, E)$  a graph:

$$\begin{aligned} n(G) = n &\geq n_0, \\ e(G) \geq \left(\frac{1}{4} + \delta\right)n^2 &\implies \# K_3 \text{'s in } G \geq cn^3 \end{aligned}$$

The above proposition “counts” the number of  $K_3$ ’s, not just showing the existence of  $K_3$  as Mantel’s theorem

**Theorem 2.7 (Mantel 1907)**

$$e(G) > \lfloor n^2/4 \rfloor \implies K_3 \subseteq G$$

## Proof outline

- Apply Szemerédi’s regularity lemma to  $G$
- Consider a regularity graph  $R$  of  $G$  (next slide)
- Conclude that  $R$  contains a copy of  $K_3$  (by Mantel’s theorem) and apply Lemma 11.2

### Whole idea

Under some assumptions, to conclude  $G$  contains many  $K_3$ , it is enough to show  $R$  contains  $K_3$

## Regularity graphs

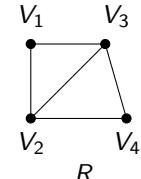
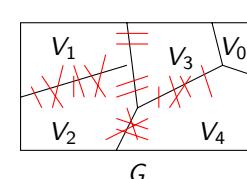
$G$  a graph;  $\varepsilon > 0$ ;

$\{V_0, V_1, \dots, V_k\}$  an  $\varepsilon$ -regular partition of  $G$ ;  $|V_1| = \dots = |V_k| = \ell$ ;  
 $0 \leq d \leq 1$

**Definition (Regularity graph)**

The **regularity graph**  $R$  of  $\{V_0, V_1, \dots, V_k\}$  with parameters  $\varepsilon, \ell, d$  is defined as follows:

- $V(R) = \{V_1, \dots, V_k\}$
- $E(R) = \{\{V_i, V_j\} \mid (V_i, V_j) \text{ an } \varepsilon\text{-regular pair of density } \geq d\}$



## Rephrasing the quantitative version

We actually prove the following

$$\forall \delta > 0 \exists n_0 \in \mathbb{N} \exists N_0 \in \mathbb{N} \exists \varepsilon > 0 \exists d \in (0, 1] \exists c > 0$$

$\forall G = (V, E)$  a graph:

$$n(G) = n \geq n_0, e(G) \geq \left(\frac{1}{4} + \delta\right) n^2$$

$\implies \forall \{V_0, V_1, \dots, V_k\}$  an  $\varepsilon$ -regular partition of  $G$

$$(n_0 \leq k \leq N_0, |V_1| = \dots = |V_k| = \ell):$$

the regularity graph of  $\{V_i\}$  with param's  $\varepsilon, \ell, d$  contains  $K_3$

$$\implies \# K_3 \text{'s in } G \geq cn^3$$

Here,  $N_0$  is a natural number that Szemerédi's regularity lemma gives

## Proof idea of Proposition 11.4 (2/11)

To show  $R$  has a lot of edges, we separately consider the following five types of edges in  $G$

- $E = E_0 \cup E_1 \cup E_2 \cup E_3 \cup E_4$
- $E_0 = \{e \in E \mid V_0 \cap e \neq \emptyset\}$
- $E_1 = \left\{ e \in E \mid \begin{array}{l} V_i \cap e \neq \emptyset, V_j \cap e \neq \emptyset \\ \text{for some } \{V_i, V_j\} \in E(R) \end{array} \right\}$
- $E_2 = \left\{ e \in E \mid \begin{array}{l} V_i \cap e \neq \emptyset, V_j \cap e \neq \emptyset \\ \text{for some } (V_i, V_j) \text{ not } \varepsilon\text{-regular} \end{array} \right\}$
- $E_3 = \left\{ e \in E \mid \begin{array}{l} V_i \cap e \neq \emptyset, V_j \cap e \neq \emptyset \\ \text{for some } (V_i, V_j) \text{ w/ } d(V_i, V_j) < d \end{array} \right\}$
- $E_4 = \{e \in E \mid e \subseteq V_i \text{ for some } V_i, i \in \{1, \dots, k\}\}$

## Proof idea of Proposition 11.4 (1/11)

Given  $\delta > 0$ , we specify  $\varepsilon = \varepsilon(\delta)$  and  $n_0 = n_0(\delta)$  later

- Apply Szemerédi's regularity lemma to  $G$  with parameters  $\varepsilon$  and  $n_0$  to obtain a number  $N_0 \in \mathbb{N}$  and an  $\varepsilon$ -regular partition  $V_0, V_1, \dots, V_k$  ( $n_0 \leq k \leq N_0$ ,  $|V_1| = \dots = |V_k| = \ell$ )
- Let  $d = d(\delta, \varepsilon, n_0)$  be constant we specify later, and  $R$  be the regularity graph of  $\{V_0, V_1, \dots, V_k\}$  with parameters  $\varepsilon, \ell, d$

## Proof idea of Proposition 11.4 (3/11)

$$E_0 = \{e \in E \mid V_0 \cap e \neq \emptyset\}$$

$$\begin{aligned} |E_0| &\leq \binom{|V_0|}{2} + |V_0||V \setminus V_0| \\ &\leq \binom{\varepsilon n}{2} + \varepsilon n \cdot k\ell \\ &\leq \frac{\varepsilon^2 n^2}{2} + \varepsilon n \cdot k \frac{n}{k} \\ &= \frac{\varepsilon^2}{2} n^2 + \varepsilon n^2 \end{aligned}$$

## Proof idea of Proposition 11.4 (4/11)

$$E_1 = \left\{ e \in E \mid \begin{array}{l} V_i \cap e \neq \emptyset, V_j \cap e \neq \emptyset \\ \text{for some } \{V_i, V_j\} \in E(R) \end{array} \right\}$$

$$\begin{aligned} |E_1| &= \sum_{\{V_i, V_j\} \in E(R)} e(V_i, V_j) \\ &\leq \sum_{\{V_i, V_j\} \in E(R)} |V_i||V_j| \\ &= \sum_{\{V_i, V_j\} \in E(R)} \ell^2 \\ &= e(R)\ell^2 \\ &\leq e(R)\left(\frac{n}{k}\right)^2 \\ &= \frac{e(R)}{k^2}n^2 \end{aligned}$$

## Proof idea of Proposition 11.4 (5/11)

$$E_2 = \left\{ e \in E \mid \begin{array}{l} V_i \cap e \neq \emptyset, V_j \cap e \neq \emptyset \\ \text{for some } (V_i, V_j) \text{ not } \varepsilon\text{-regular} \end{array} \right\}$$

$$\begin{aligned} |E_2| &\leq (\# \text{ non } \varepsilon\text{-regular pairs}) \cdot \ell^2 \\ &\leq \varepsilon k^2 \ell^2 \\ &\leq \varepsilon k^2 \left(\frac{n}{k}\right)^2 \\ &= \varepsilon n^2 \end{aligned}$$

Recall

# non  $\varepsilon$ -regular pairs in  $V_1, \dots, V_k \leq \varepsilon k^2$  in any  $\varepsilon$ -regular partition

## Proof idea of Proposition 11.4 (6/11)

$$E_3 = \left\{ e \in E \mid \begin{array}{l} V_i \cap e \neq \emptyset, V_j \cap e \neq \emptyset \\ \text{for some } (V_i, V_j) \text{ w/ } d(V_i, V_j) < d \end{array} \right\}$$

$$\begin{aligned} |E_3| &\leq (\# \text{ pairs}) \cdot d\ell^2 \\ &= \binom{k}{2} \cdot d\ell^2 \leq \frac{k^2}{2} \cdot d \cdot \left(\frac{n}{k}\right)^2 = \frac{d}{2}n^2 \end{aligned}$$

$$E_4 = \{ e \in E \mid e \subseteq V_i \text{ for some } V_i, i \in \{1, \dots, k\} \}$$

$$|E_4| \leq k \cdot \binom{\ell}{2} \leq k \cdot \frac{\ell^2}{2} \leq \frac{k}{2} \left(\frac{n}{k}\right)^2 = \frac{1}{2k}n^2$$

## Proof idea of Proposition 11.4 (7/11)

- Therefore,

$$\begin{aligned} |E| &\leq |E_0| + |E_1| + |E_2| + |E_3| + |E_4| \\ &\leq \left(\frac{\varepsilon^2}{2} + \varepsilon\right)n^2 + \frac{e(R)}{k^2}n^2 + \varepsilon n^2 + \frac{d}{2}n^2 + \frac{1}{2k}n^2 \\ &\leq \left(\frac{e(R)}{k^2} + \frac{1}{2}\varepsilon^2 + 2\varepsilon + \frac{d}{2} + \frac{1}{2k}\right)n^2 \end{aligned}$$

- On the other hand, from the assumption we have

$$|E| \geq \left(\frac{1}{4} + \delta\right)n^2$$

## Proof idea of Proposition 11.4 (8/11)

- Therefore,

$$\begin{aligned}\frac{1}{4} + \delta &\leq \frac{e(R)}{k^2} + \frac{1}{2}\varepsilon^2 + 2\varepsilon + \frac{d}{2} + \frac{1}{2k} \\ \therefore e(R) &\geq \left(\frac{1}{4} + \delta - \frac{1}{2}\varepsilon^2 - 2\varepsilon - \frac{d}{2} - \frac{1}{2k}\right) k^2\end{aligned}$$

- Setting  $\varepsilon = \frac{\delta}{16}$ ,  $n_0 = \left\lceil \frac{1}{\delta} \right\rceil$ ,  $d = \frac{\delta}{4}$  we have

$$e(R) \geq \left(\frac{1}{4} + \gamma\right) k^2$$

for some  $\gamma$  only depending on  $\delta$  (calculation in the next slide)

## Proof idea of Proposition 11.4 (10/11)

We now have

$$e(R) \geq \left(\frac{1}{4} + \gamma\right) k^2$$

for some  $\gamma$  only depending on  $\delta$ .

- $\therefore R$  contains a copy of  $K_3$  (from Mantel's theorem)
- $\therefore$  The  $\varepsilon$ -regular partition of  $G$  has three parts, say  $V_1, V_2, V_3$ , such that  $(V_i, V_j)$  are  $\varepsilon$ -regular ( $1 \leq i < j \leq 3$ ) with density  $d$
- $\therefore V_1, V_2, V_3$  contain at least  $(1-4\varepsilon)(d-\varepsilon)^3\ell^3$   $K_3$ 's  
(by Lemma 11.2 and  $\varepsilon = \delta/16 < \delta/8 = d/2$ )
- Namely,  $G$  contains  $(1-4\varepsilon)(d-\varepsilon)^3\ell^3$   $K_3$ 's

## Proof idea of Proposition 11.4 (9/11)

$$\begin{aligned}\frac{1}{4} + \delta - \frac{1}{2}\varepsilon^2 - 2\varepsilon - \frac{d}{2} - \frac{1}{2k} &\geq \frac{1}{4} + \delta - \frac{1}{512}\delta^2 - \frac{1}{8}\delta - \frac{1}{8}\delta - \frac{1}{2}\delta && (k \geq n_0 \geq \frac{1}{\delta}) \\ &= \frac{1}{4} + \frac{1}{4}\delta - \frac{1}{512}\delta^2 \\ &= \frac{1}{4} + \frac{\delta}{512}(128 - \delta) \\ &\geq \frac{1}{4} + \frac{\delta}{512}(128 - 1) && (\because \text{we may assume } \delta \leq 1) \\ &= \frac{1}{4} + \frac{127}{512}\delta; \\ \text{Then set } \gamma &= \frac{127}{512}\delta\end{aligned}$$

## Proof idea of Proposition 11.4 (11/11)

How many  $K_3$ 's at least in  $G$ , eventually?

$$\begin{aligned}(1-4\varepsilon)(d-\varepsilon)^3\ell^3 &\geq (1-4\varepsilon)(d-\varepsilon)^3 \left(\frac{(1-\varepsilon)n}{k}\right)^3 \\ &= (1-4\varepsilon)(d-\varepsilon)^3(1-\varepsilon)^3 \frac{n^3}{k^3} \\ &\geq (1-4\varepsilon)(d-\varepsilon)^3(1-\varepsilon)^3 \frac{n^3}{N_0^3} && (N_0 \text{ in Regularity lemma}) \\ &= c(\varepsilon, d, N_0)n^3\end{aligned}$$

and  $\varepsilon, d$  only depend on  $\delta$ ;  $N_0$  only depends on  $\varepsilon, n_0$  which only depend on  $\delta$ ;  $\therefore c$  only depends on  $\delta$   $\square$

## Quantitative version of Turán's theorem

## Lemma 11.5 (Quantitative version of Turán's theorem)

$\forall r \in \mathbb{N} \forall \delta > 0 \exists n_0 \in \mathbb{N} \exists N_0 \in \mathbb{N} \exists \varepsilon > 0 \exists d \in (0, 1] \exists c > 0$   
 $\forall G = (V, E) \text{ a graph:}$

$$n(G) = n \geq n_0, e(G) \geq \left(1 + \frac{1}{r-1} + \delta\right) \binom{n}{2}$$

$\implies \forall \{V_0, V_1, \dots, V_k\}$  an  $\varepsilon$ -regular partition of  $G$   
 $(n_0 \leq k \leq N_0, |V_1| = \dots = |V_k| = \ell)$ :

the regularity graph of  $\{V_i\}$  with param's  $\varepsilon, \ell, d$  contains  $K_r$

$\implies \# K_r \text{'s in } G \geq cn^r$

Here,  $N_0$  is a natural number that Szemerédi's regularity lemma gives

Proof

Exercise

## Today's contents

① Review of Lecture 9

② Density and  $\varepsilon$ -regular pairs

③ Szemerédi's regularity lemma

④ Embedding lemma and a proof of the Erdős-Stone theorem

⑤ Algorithmic issue

I think we somewhat missed our road map...

We want to show the following

## Rephrasing the Erdős-Stone thm

$\forall s \geq 1, r \geq 2$  natural numbers  $\forall \delta > 0 \exists n_1 \in \mathbb{N} \forall n \geq n_1 \forall G:$

$n(G) = n \geq n_1$  and

$$e(G) \geq \left(1 - \frac{1}{r-1} + \delta\right) \binom{n}{2} \Rightarrow G \supseteq T_{rs,r}$$

Proof outline

- Use Szemerédi's regularity lemma to prove the embedding lemma
- Use the embedding lemma to prove Erdős-Stone thm

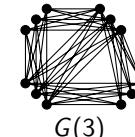
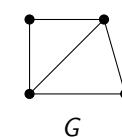
## Expanding vertices in a graph

$G = (V, E)$  a graph;  $s \in \mathbb{N}$

## Notation

The graph  $G(s)$  is defined as follows:

- $V(G(s)) = V \times \{1, \dots, s\}$
- $E(G(s)) = \{\{(u, i), (v, j)\} \mid \{u, v\} \in E\}$



Note:  $K_r(s) \simeq T_{rs,r}$

## Embedding lemma

### Embedding lemma (Komlós, Simonovits '96)

- $\forall d \in (0, 1] \forall \Delta \geq 1$
- $\exists \gamma_0 > 0$
- $\forall G$  a graph  $\forall H$  a graph w/  $\Delta(H) \leq \Delta$
- $\forall \gamma < \gamma_0 \forall s \in \mathbb{N} \forall \ell \geq s/\gamma_0$
- $\forall \{V_0, V_1, \dots, V_k\}$  a  $\gamma$ -reg partition of  $G$  ( $|V_1| = \dots = |V_k| = \ell$ ):  
 $R$  is the reg graph of  $\{V_i\}$   
w/ param's  $\gamma, \ell, d$ , and  $\implies H \subseteq G$   
 $H \subseteq R(s)$

Under some assumptions, to conclude  $G$  contains  $H$ , instead of showing  $R$  contains  $H$ , it is enough to show  $R(s)$  contains  $H$

## Proof idea of the Erdős-Stone thm (1/4)

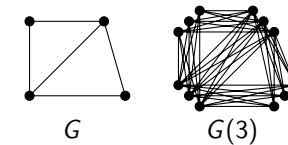
### Rephrasing the Erdős-Stone thm

$\forall s \geq 1, r \geq 2$  natural numbers  $\forall \delta > 0 \exists n_1 \in \mathbb{N} \forall n \geq n_1 \forall G: n(G) = n \geq n_1$  and

$$e(G) \geq \left(1 - \frac{1}{r-1} + \delta\right) \binom{n}{2} \Rightarrow G \supseteq T_{rs,r}$$

### Proof outline:

- Enough to show  $T_{rs,r} \subseteq R(s)$  under the given assumption (by the embedding lemma)
- $\therefore$  enough to show  $K_r \subseteq R$  under the assumptions



## Proof idea of the Erdős-Stone thm (2/4)

We use the following as well

### Lemma 11.5 (Quantitative version of Turán's theorem)

$\forall r \in \mathbb{N} \forall \delta > 0 \exists n_2 \in \mathbb{N} \exists N_2 \in \mathbb{N} \exists \varepsilon_2 > 0 \exists d_2 \in (0, 1]$

$\forall G = (V, E)$  a graph:

$$n(G) = n \geq n_2, e(G) \geq \left(1 + \frac{1}{r-1} + \delta\right) \binom{n}{2}$$

$\implies \forall \{V_0, V_1, \dots, V_k\}$  an  $\varepsilon_2$ -regular partition of  $G$

( $n_2 \leq k \leq N_2, |V_1| = \dots = |V_k| = \ell$ ):

the regularity graph of  $\{V_i\}$  with param's  $\varepsilon_2, \ell, d_2$  contains  $K_r$

Here,  $N_2$  is a natural number that Szemerédi's regularity lemma gives

## Proof idea of the Erdős-Stone thm (3/4)

Given  $s \geq 1, r \geq 2, \delta > 0$  (as in the Erdős-Stone thm)

- Lem 11.5 gives  $n_2(r, \delta), N_2(r, \delta), \varepsilon_2(r, \delta), d_2(r, \delta)$

- Let  $d = d_2$

- Let  $\Delta = \Delta(T_{rs,r}) = (r-1)s$

- Embedding Lemma gives  $\gamma_0 = \gamma_0(d, \Delta)$

- Let  $\varepsilon < \min\{\gamma_0, \varepsilon_2\}$

(Note:  $\varepsilon < \varepsilon_2$ )

- Let  $n_1 \geq \max \left\{ n_2, \frac{N_2 s}{\gamma_0(1-\varepsilon)} \right\}$

## Proof idea of the Erdős-Stone thm (4/4)

Let  $G$  satisfy  $n(G) = n \geq n_1$ , and  $H$  satisfy  $\Delta(H) \leq \Delta$   
 (Note:  $n \geq n_2$ )

- By Regularity Lemma,  $\exists$  an  $\varepsilon$ -reg partition  $\{V_0, V_1, \dots, V_k\}$  w/  
 $|V_1| = \dots = |V_k| = \ell$  and  $n_2 \leq k \leq N_2$
- Note: this  $\varepsilon$ -reg partition is  $\varepsilon_2$ -reg as well
- Let  $R$  be the regularity graph of  $\{V_i\}$  w/ parameters  $\varepsilon, \ell, d$
- By Lem 11.5 (and  $\varepsilon < \varepsilon_2$ ,  $d \leq d_2$ ),  $K_r \subseteq R$
- $\therefore T_{rs,r} \subseteq R(s)$
- Let  $\gamma = \varepsilon$  (Note  $\varepsilon < \min\{\gamma_0, \varepsilon_2\} \leq \gamma_0$ )
- Then  $\ell \geq \frac{(1-\varepsilon)n}{N_0} \geq \frac{(1-\varepsilon)n_1}{N_0} \geq \frac{1-\varepsilon}{N_0} \frac{N_0 s}{\gamma_0(1-\varepsilon)} = \frac{s}{\gamma_0}$
- $\therefore$  By Embedding Lemma,  $T_{rs,r} \subseteq G$  □

## Today's contents

- ① Review of Lecture 9
- ② Density and  $\varepsilon$ -regular pairs
- ③ Szemerédi's regularity lemma
- ④ Embedding lemma and a proof of the Erdős-Stone theorem
- ⑤ Algorithmic issue

Algorithmic issue  
Deciding a given pair is  $\varepsilon$ -regularProblem  $\varepsilon$ -REGULARITY

Input: A graph  $G = (V, E)$  with its partite sets  $A, B$  and a real number  $\varepsilon > 0$

Question: Is  $\{A, B\}$   $\varepsilon$ -regular?

Fact (Alon, Duke, Lefmann, Rödl, Yuster '94)

Problem  $\varepsilon$ -REGULAR is coNP-complete

Algorithmic issue  
Finding a  $\varepsilon$ -regular partition

Now comes a strange phenomenon... What if we want to construct one partition?

Problem  $\varepsilon$ -REGULARITY PARTITION

Input: A graph  $G = (V, E)$  and a real number  $\varepsilon > 0$

Output: An  $\varepsilon$ -regular partition of  $G$

Fact (Alon, Duke, Lefmann, Rödl, Yuster '94)

Problem  $\varepsilon$ -REGULAR PARTITION can be solved in polynomial time

This “algorithmic version” of Szemerédi’s regularity lemma is quite useful (in theory)

### Algorithmic issue Colorable? or far from being colorable?

The following was proved with the regularity lemma

#### Fact (Duke, Rödl '85)

$\forall \varepsilon > 0 \ \forall k \in \mathbb{N} \ \exists c \ \exists \delta \ \forall G$  an  $n$ -vertex graph:

$G$  contains a  $c$ -vertex non- $k$ -colorable subgraph or  
 $G$  has  $\delta n^2$  edges whose removal leaves a  $k$ -colorable graph

The following is an algorithmic version

#### Fact (Alon, Duke, Lefmann, Rödl, Yuster '94)

Fix  $\varepsilon > 0$ ,  $k \in \mathbb{N}$ . For a given  $n$ -vertex graph  $G$  we can find either

an  $O(1)$ -vertex non- $k$ -colorable subgraph of  $G$  or  
 $O(n^2)$  edges whose removal leaves a  $k$ -colorable graph

in polynomial time

### Algorithmic issue

#### Approximately finding a maximum edge cut

#### Fact (Frieze, Kannan '96)

Fix  $\varepsilon > 0$ . Given an  $n$ -vertex graph  $G$ , we can compute in polynomial time an edge cut  $[X, \bar{X}]$  of  $G$  that satisfies

$$|[X^*, \bar{X}^*]| - |[X, \bar{X}]| \leq \varepsilon n^2$$

where  $[X^*, \bar{X}^*]$  is an edge cut with maximum number of edges

- Their algorithm uses ideas from Szemerédi's regularity lemma (Another algorithm by Arora, Karger, Karpinski '95 is based on integer linear programming and random sampling)
- If  $G$  is "dense" ( $|[X^*, \bar{X}^*]| = \Omega(n^2)$ ) then this algorithm is a polynomial-time approximation scheme (PTAS)
- The " $\ell$ -way version" is also available (Frieze, Kannan '96)