Linear-Time Counting Algorithms for Independent Sets in Chordal Graphs

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University of Metz, Metz, France

$\boxed{\text{Setup}}$ $G = (V, E)$ a graph (undirected, finite, simple)

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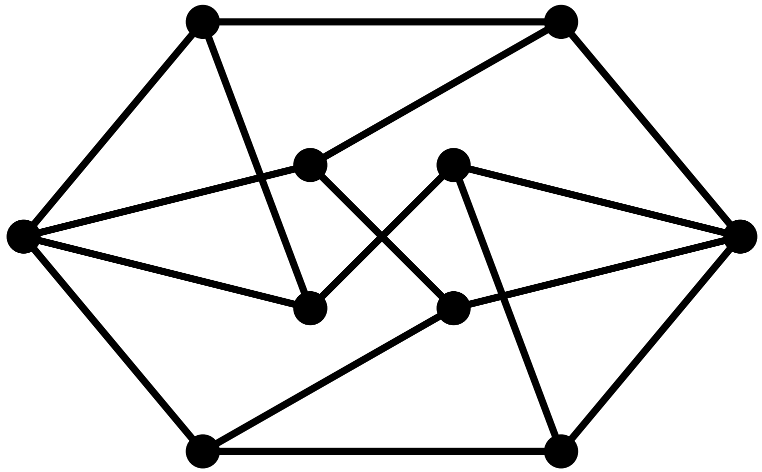
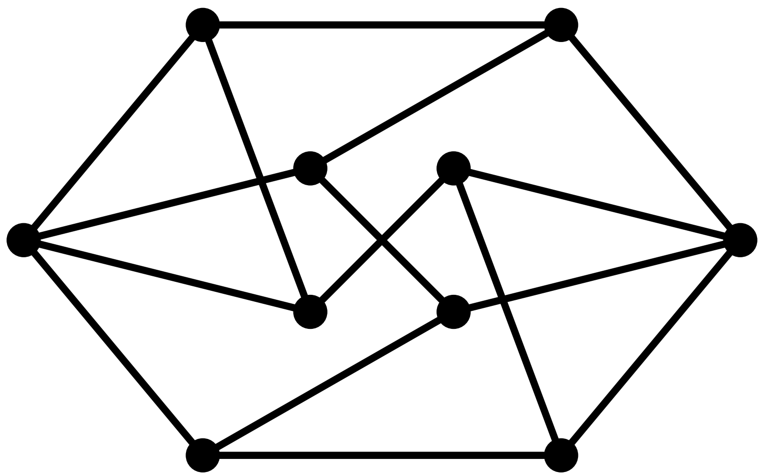


Setup $G = (V, E)$ a graph (undirected, finite, simple)

Def. A set $I \subseteq V$ is an **independent set** of G if no two vertices in I are adjacent.



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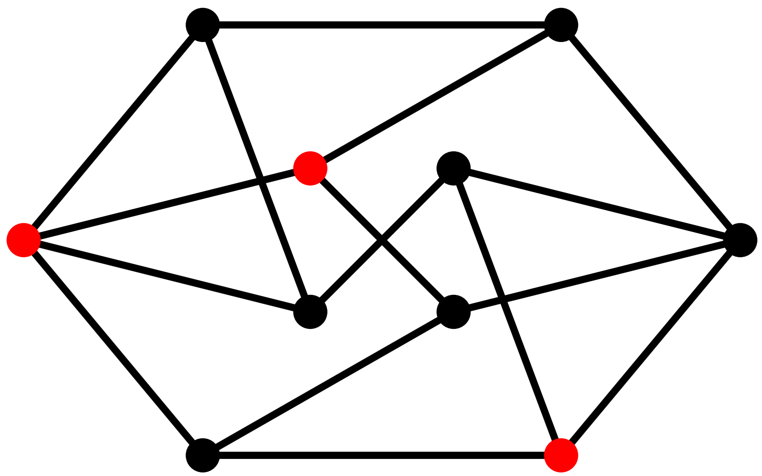

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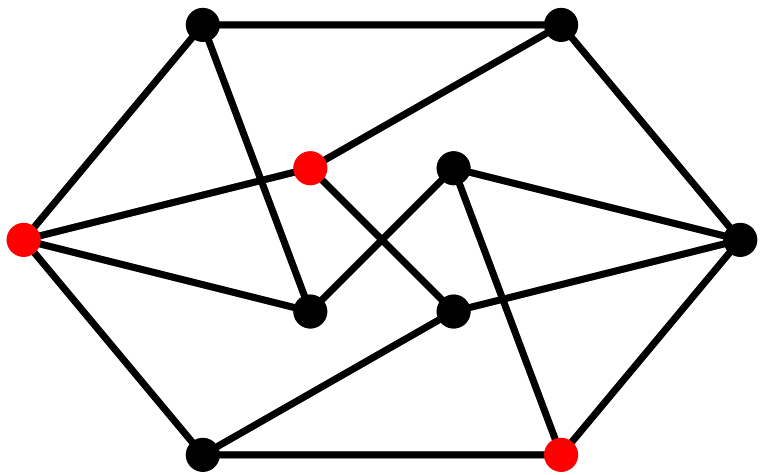

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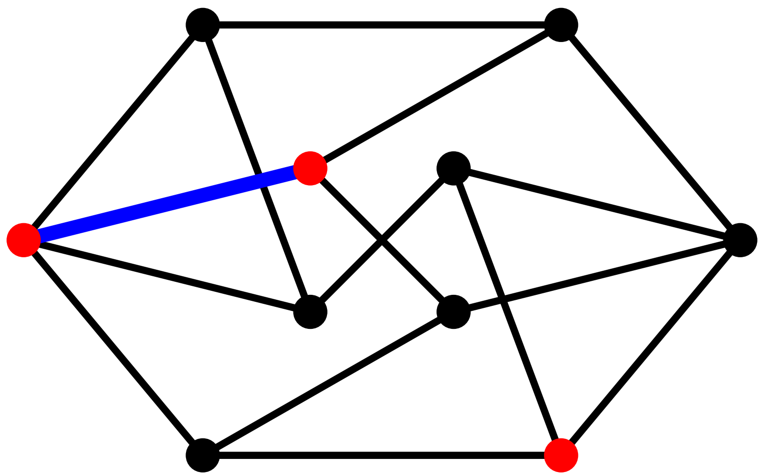

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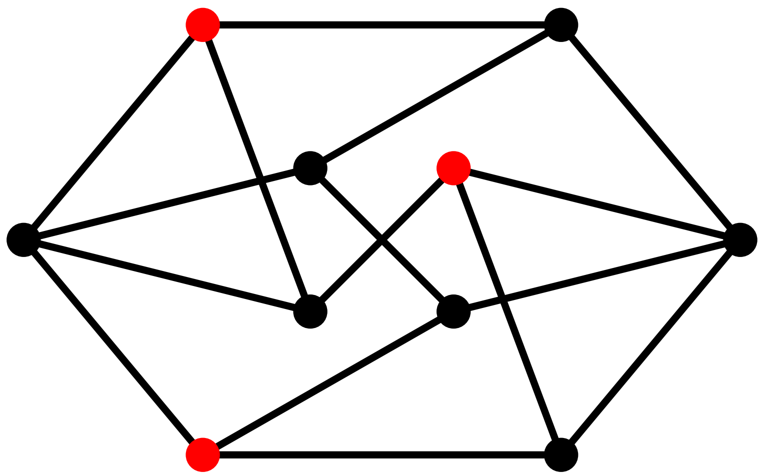

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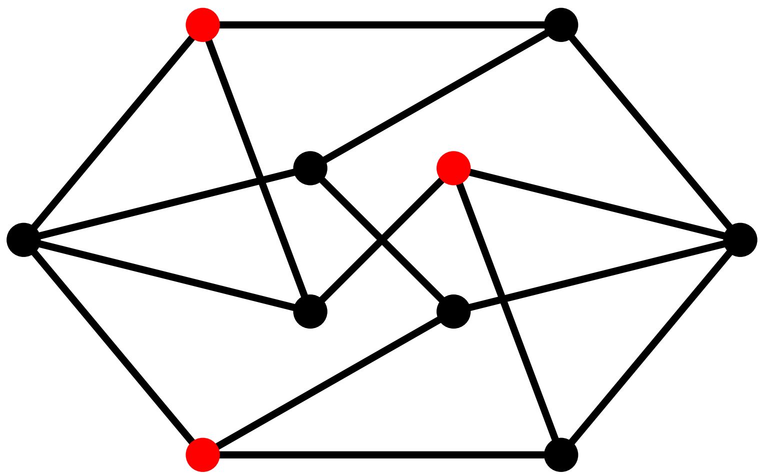

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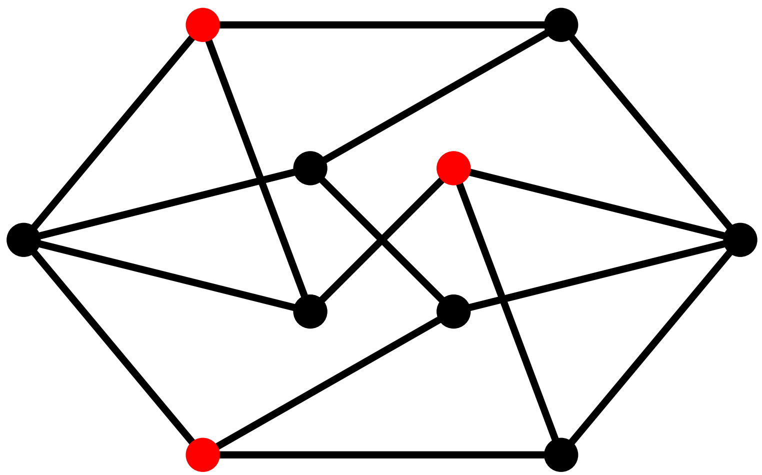
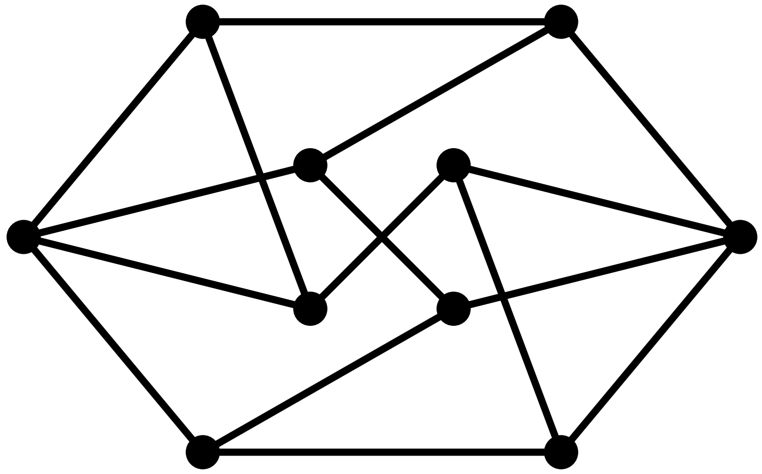


An independent set

Also called a stable set of G

$\boxed{\text{Input}}$ $G = (V, E)$ a graph

$\boxed{\text{Output}}$ All independent sets of G



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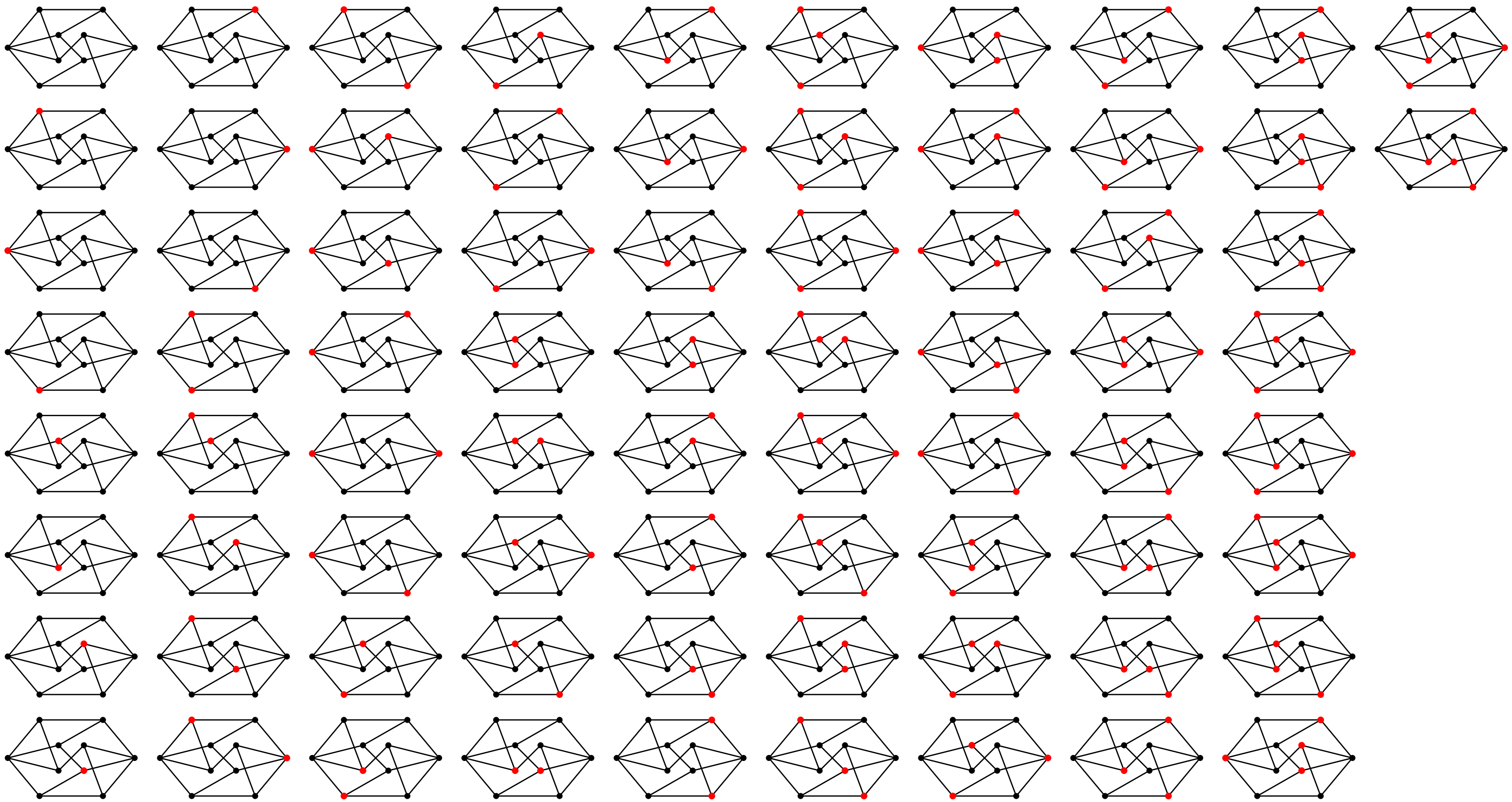
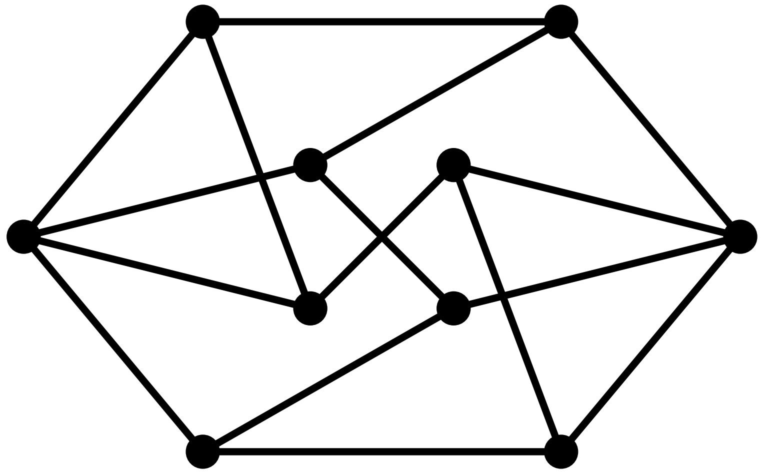


74 independent sets in total

Input $G = (V, E)$ a graph

Output # independent sets of G



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Output # independent sets of G

74

We study the following **counting problems**:

Input $G = (V, E)$ a graph

Output (1) # independent sets of G

(2) # maximum independent sets of G

(3) # independent sets of G of fixed size

(4) # maximal independent sets of G

(5) # minimum maximal independent sets of G

A maximum independent set

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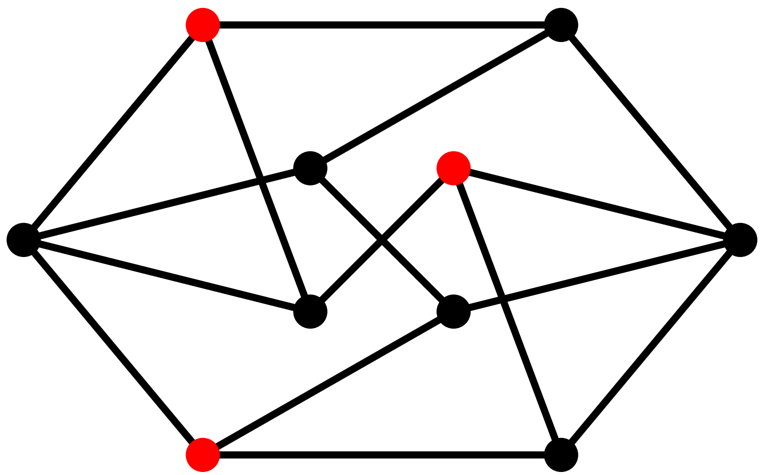

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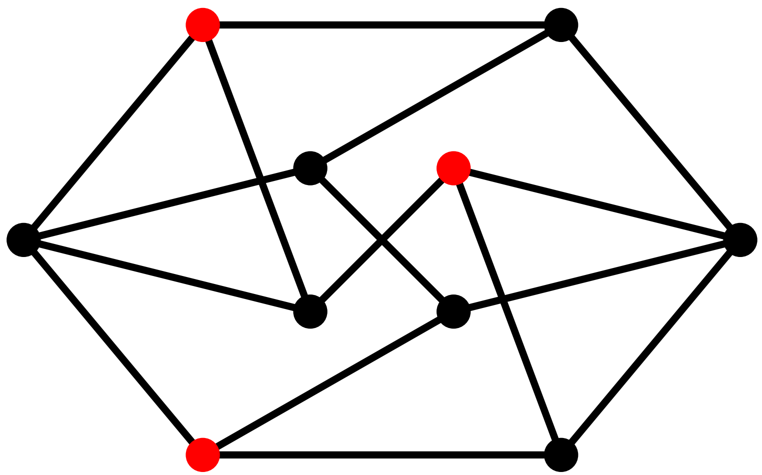

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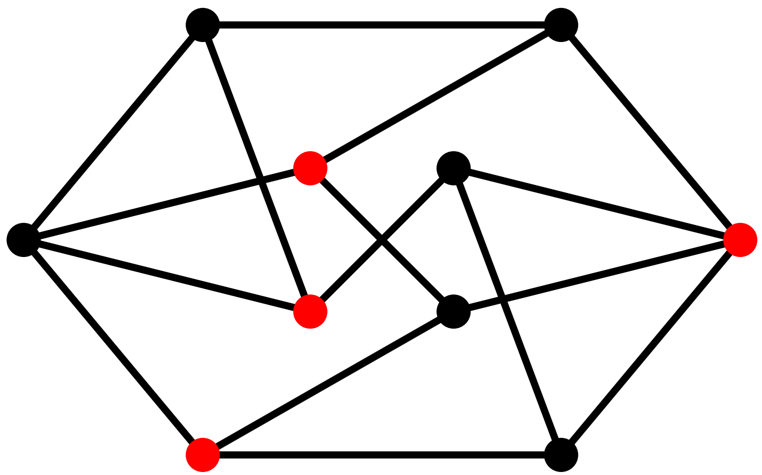

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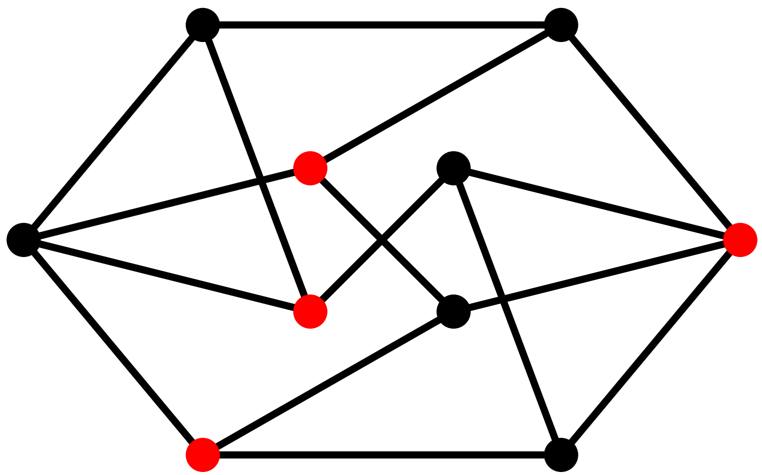
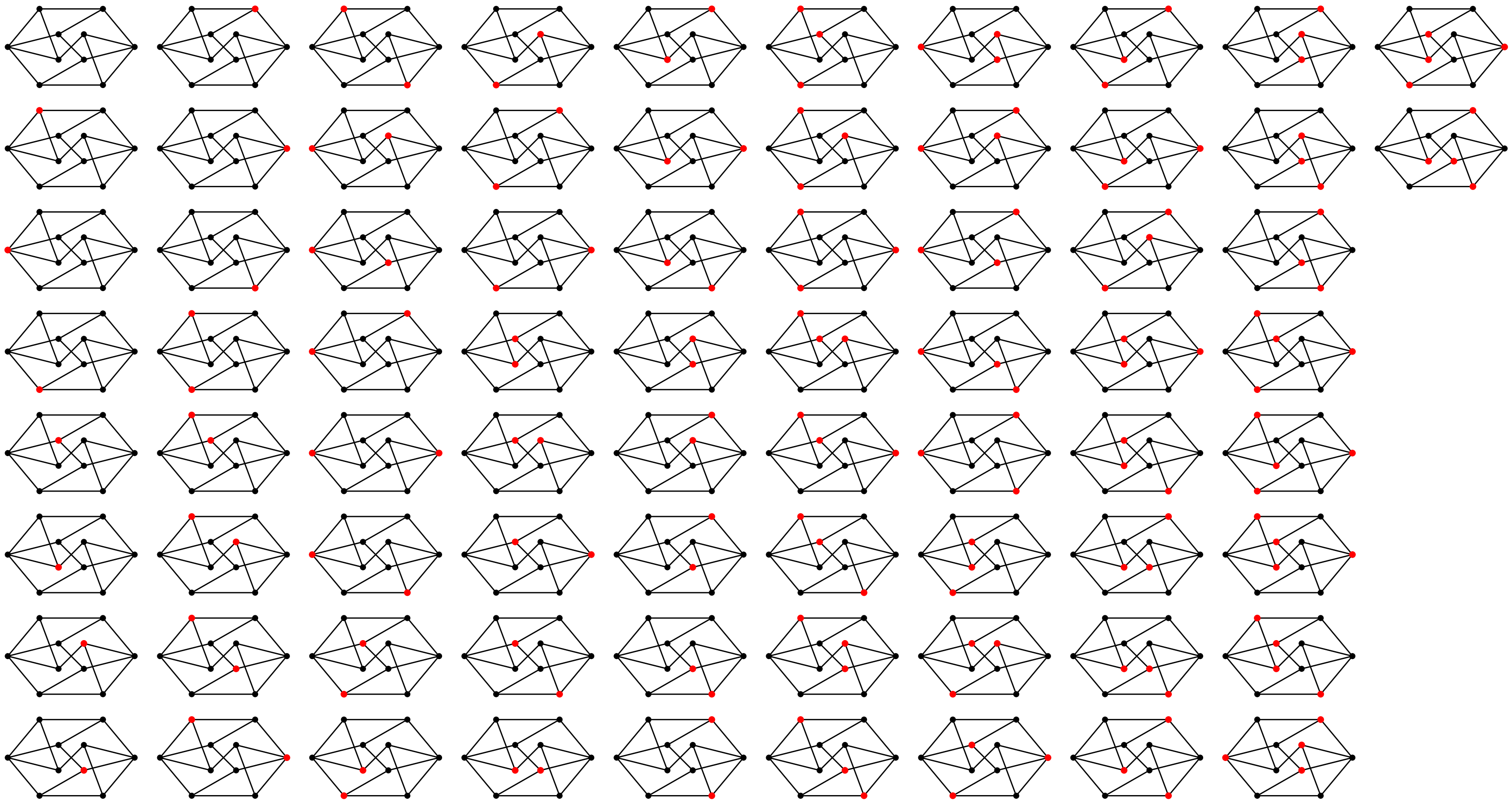


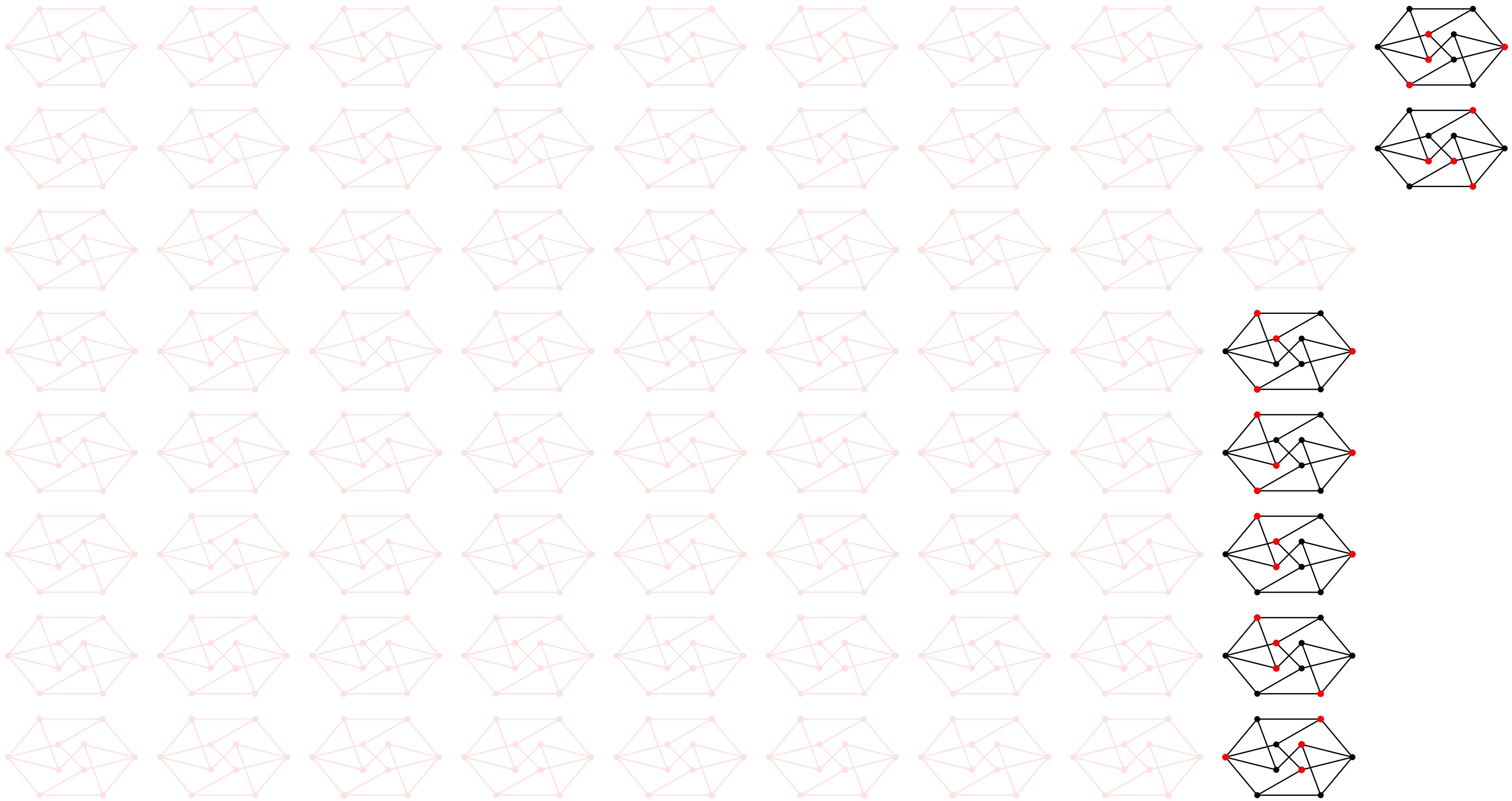
Maximum

Maximum independent sets



74 independent sets

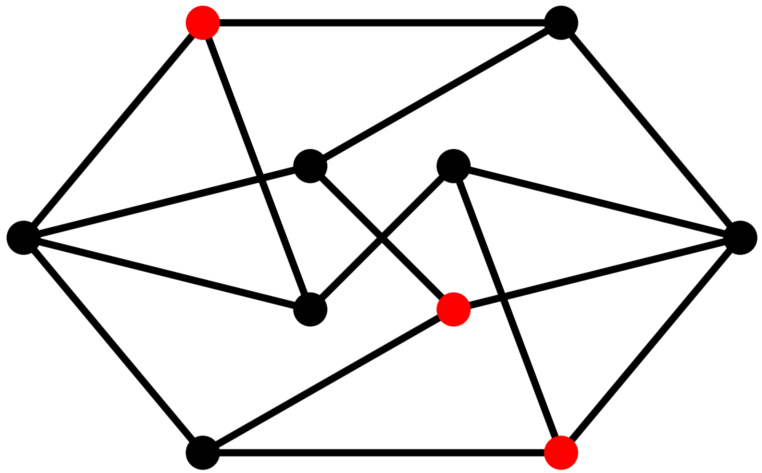
Maximum independent sets



74 independent sets 7 maximum independent sets

Setup $G = (V, E)$ a graph

Def. An independent set I of G is **maximal** if no proper superset of I is independent.



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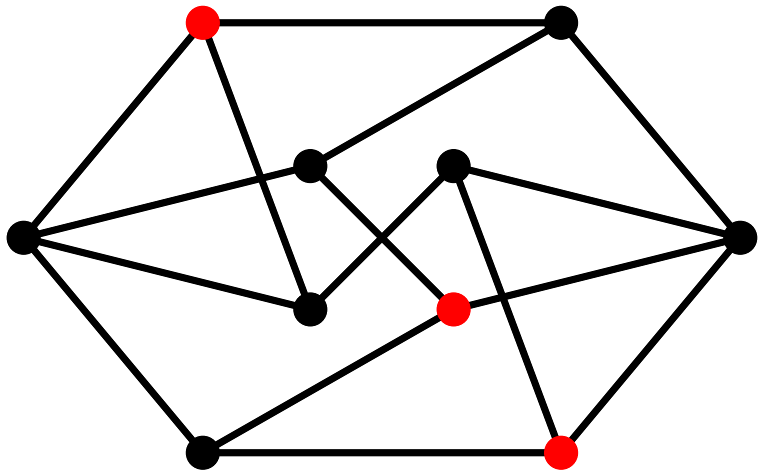

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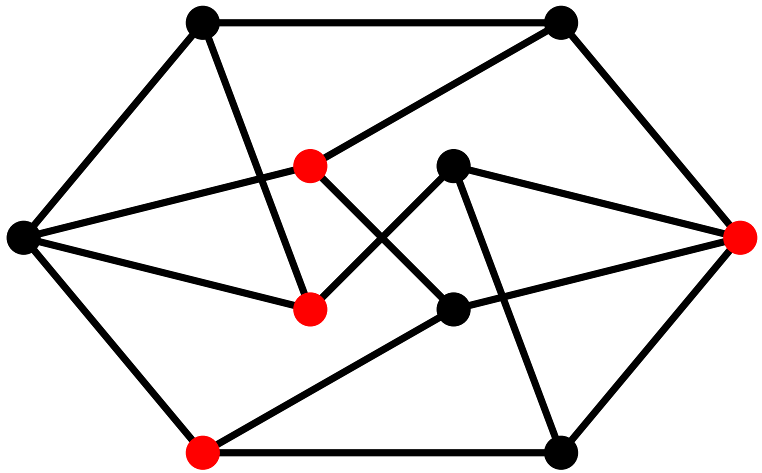

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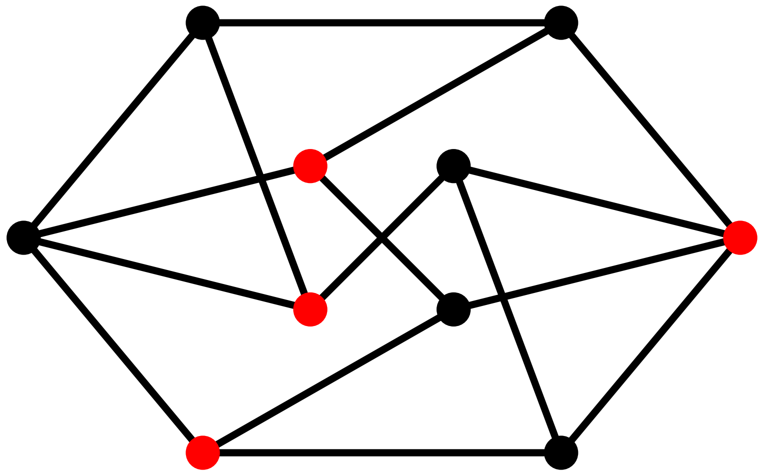

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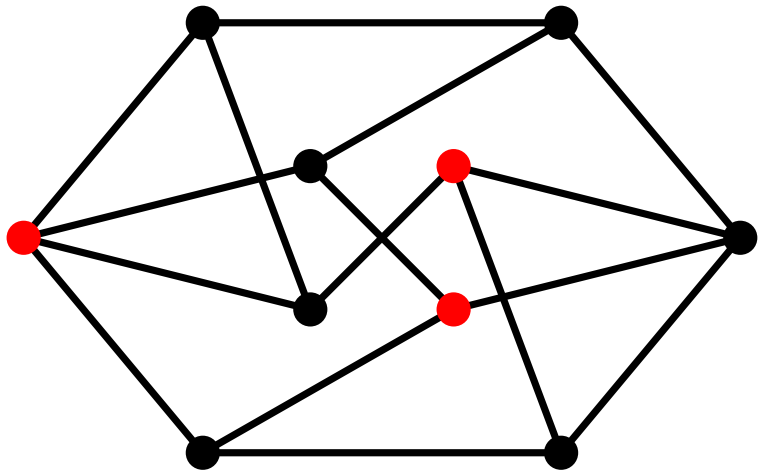

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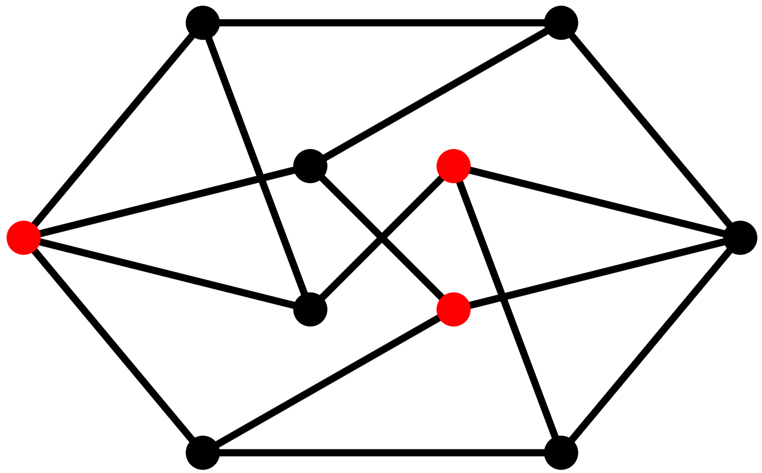

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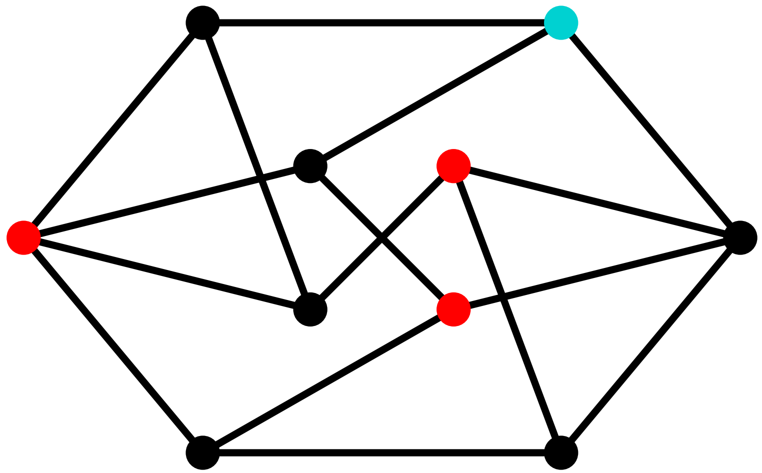

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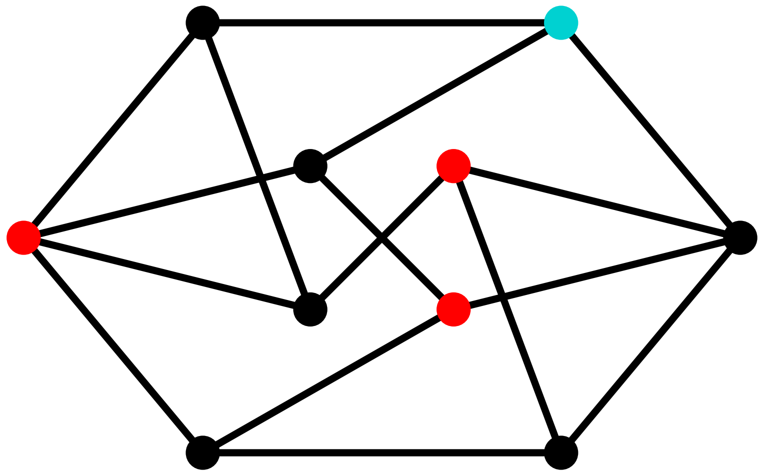
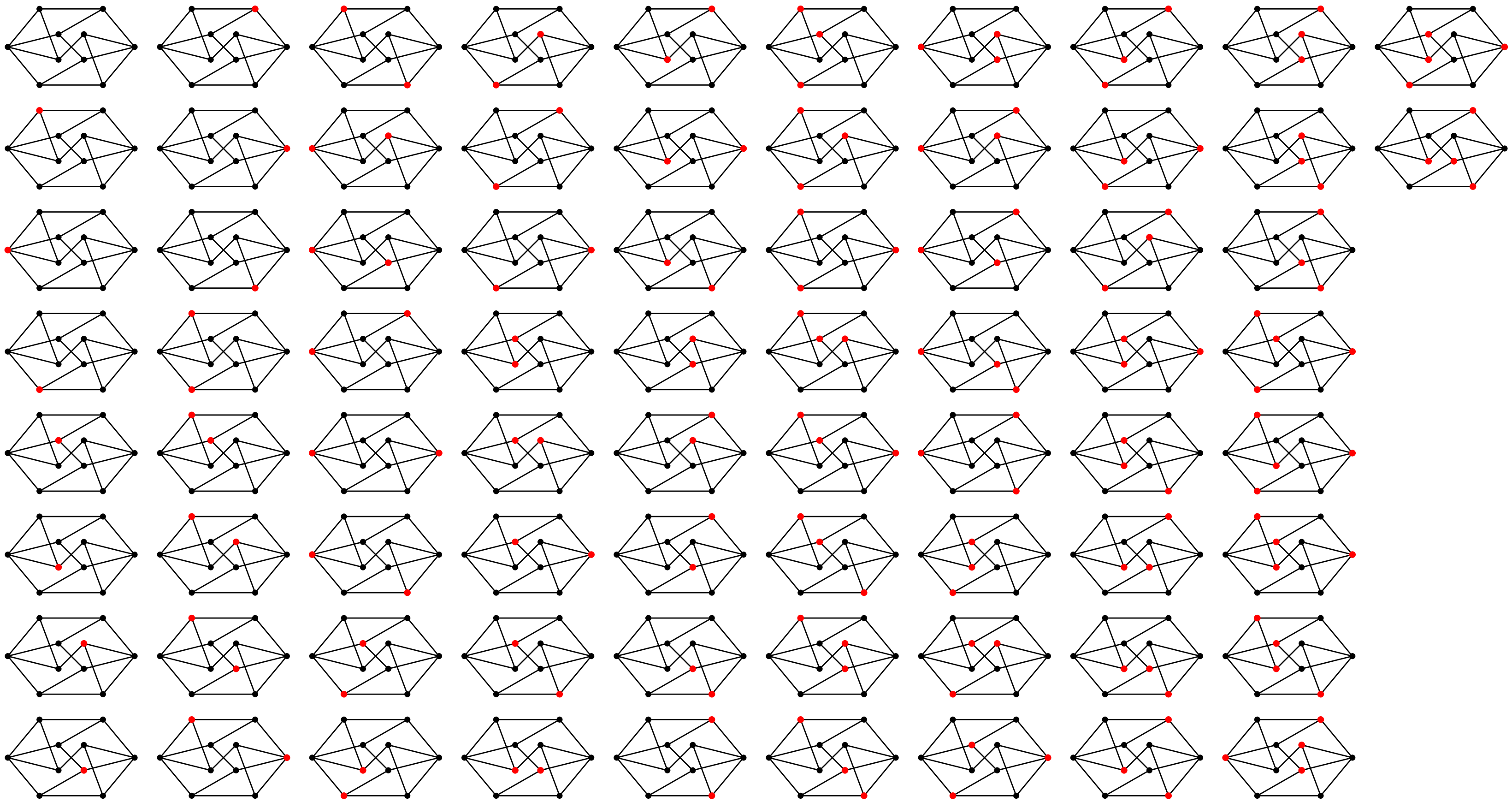


Not maximal

Also called an independent dominating set of G

Maximal independent sets



74 independent sets

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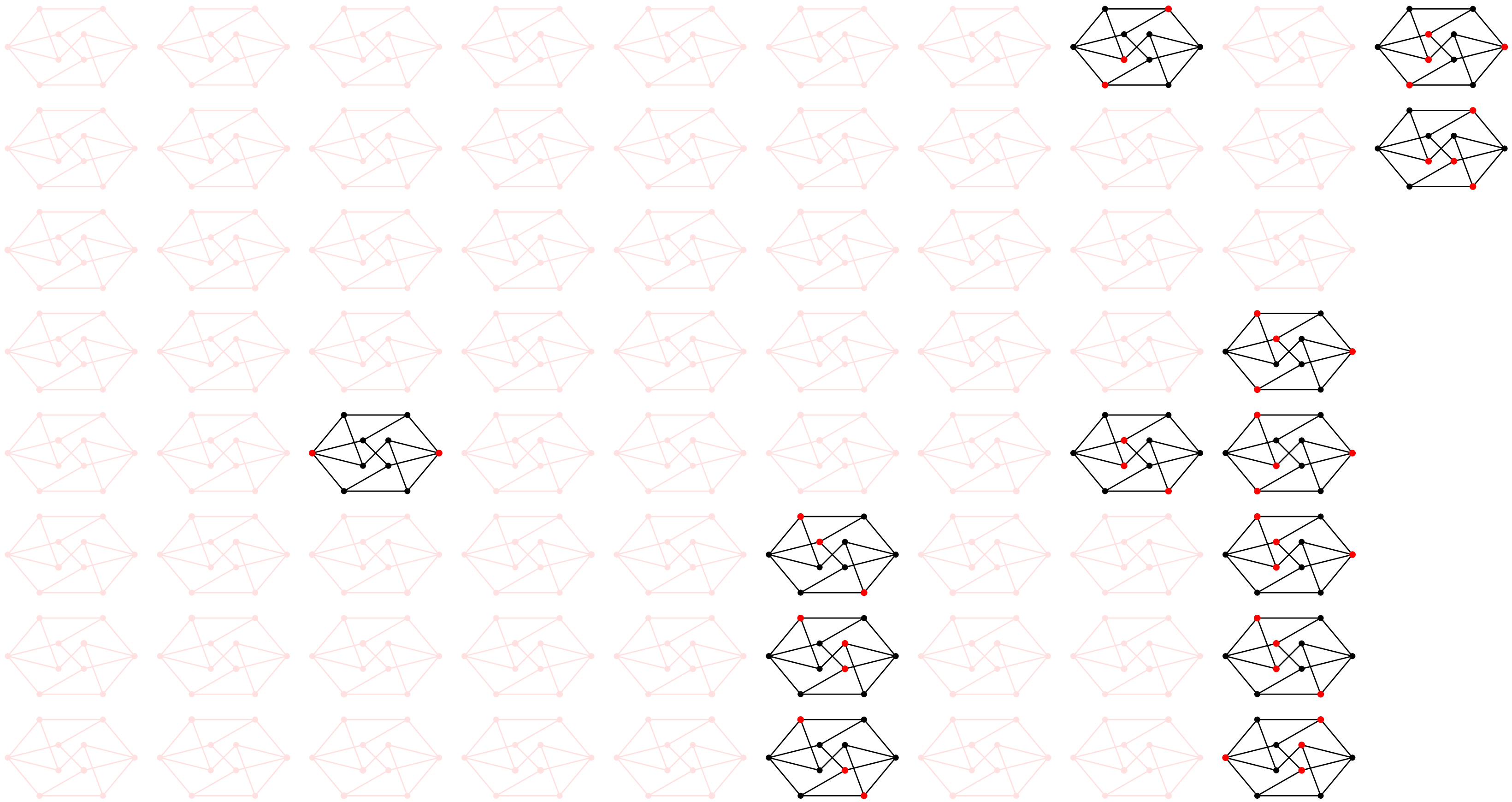
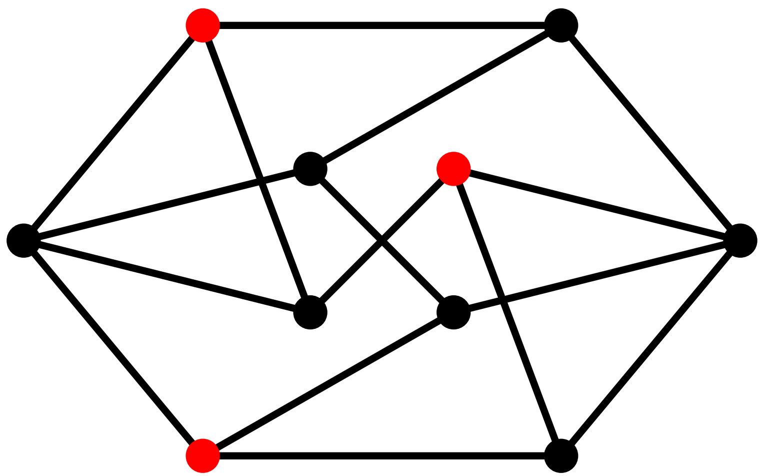


74 independent sets 13 maximal independent sets

Setup $G = (V, E)$ a graph

Def. An independent set I of G is **minimum maximal** if it is maximal and has the smallest size among all maximal independent sets of G.



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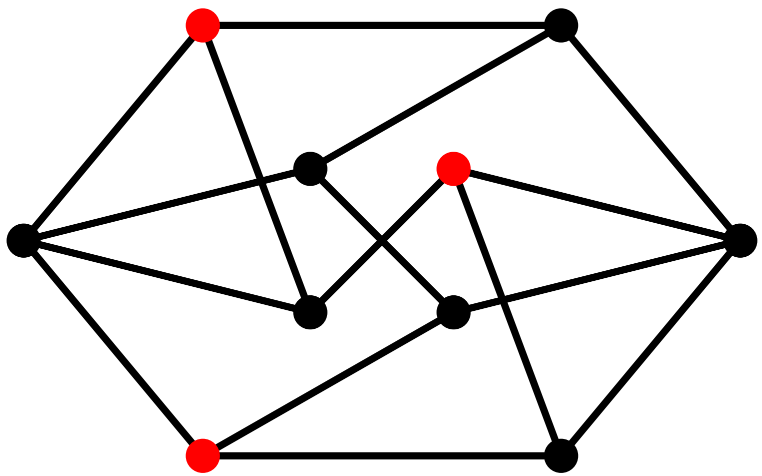

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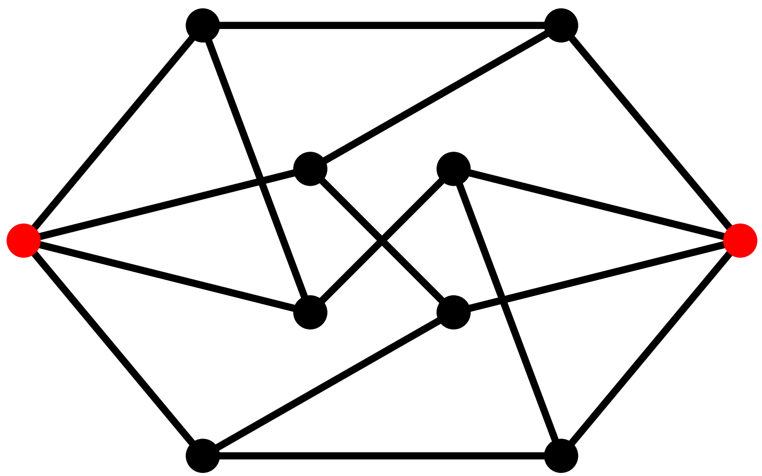

A minimum maximal independent set

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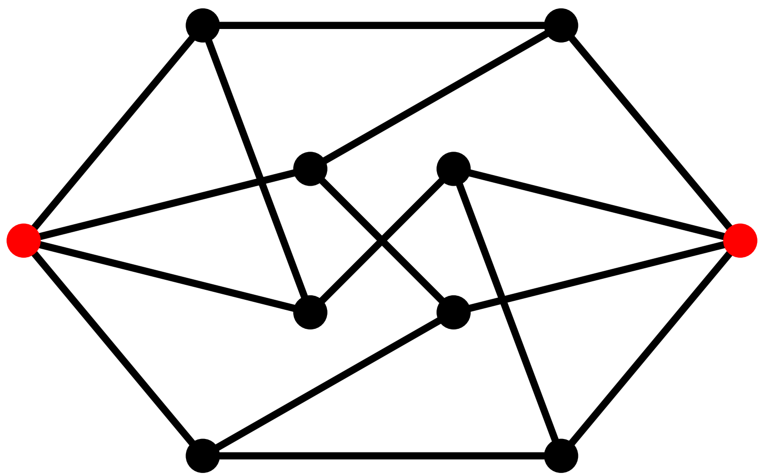

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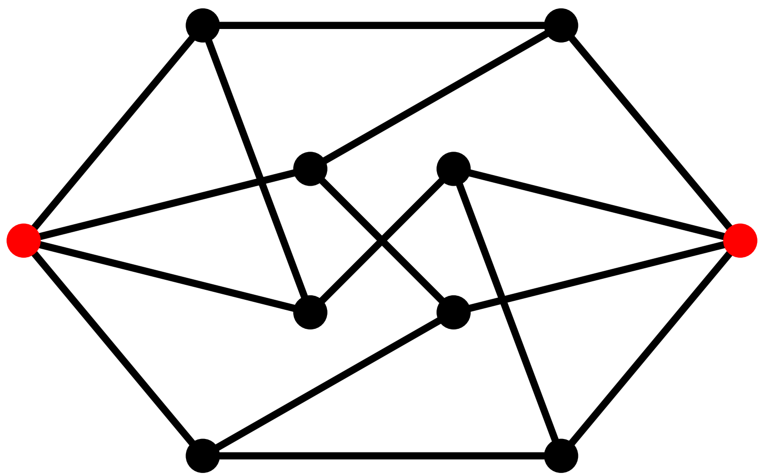
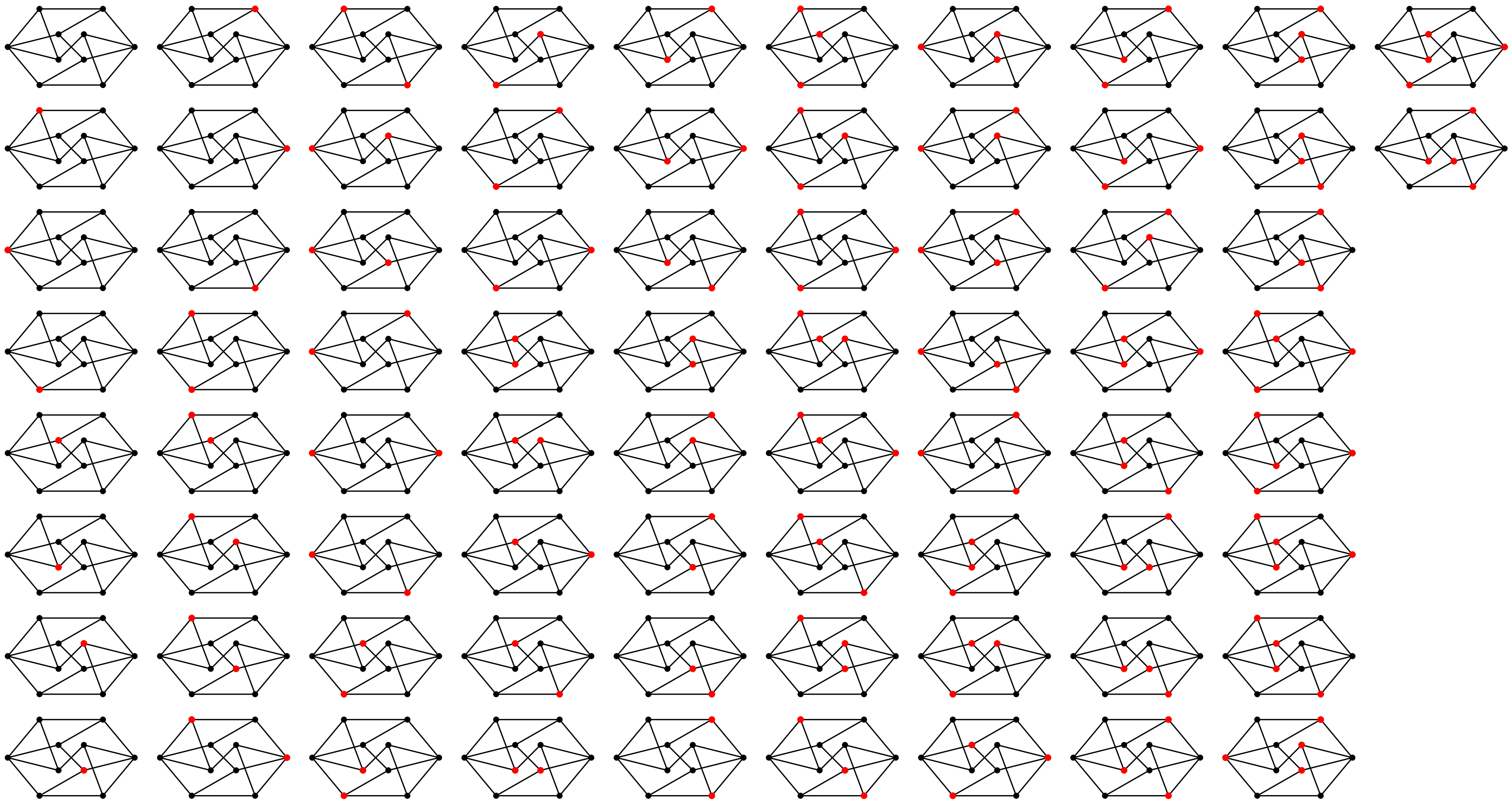


Minimum maximal

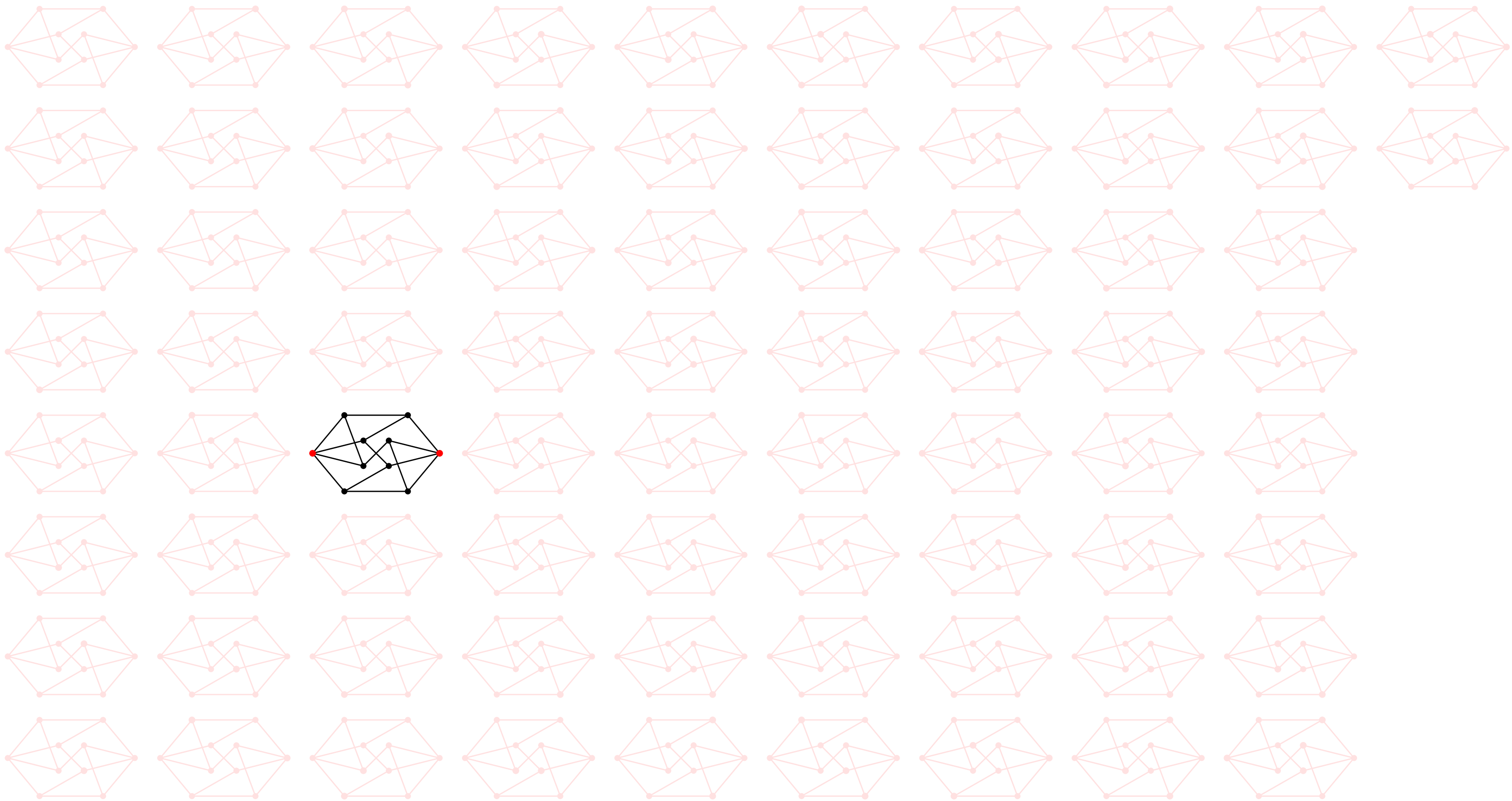
Also called a minimum independent dominating set of G

Minimum maximal independent sets



74 independent sets

Minimum maximal independent sets



74 independent sets

1 minimum maximal independent set

We study the following **counting problems**:

Input $G = (V, E)$ a graph

Output

(1) # independent sets of G **74**

(2) # maximum independent sets of G **7**

(3) # independent sets of G of fixed size

(4) # maximal independent sets of G **13**

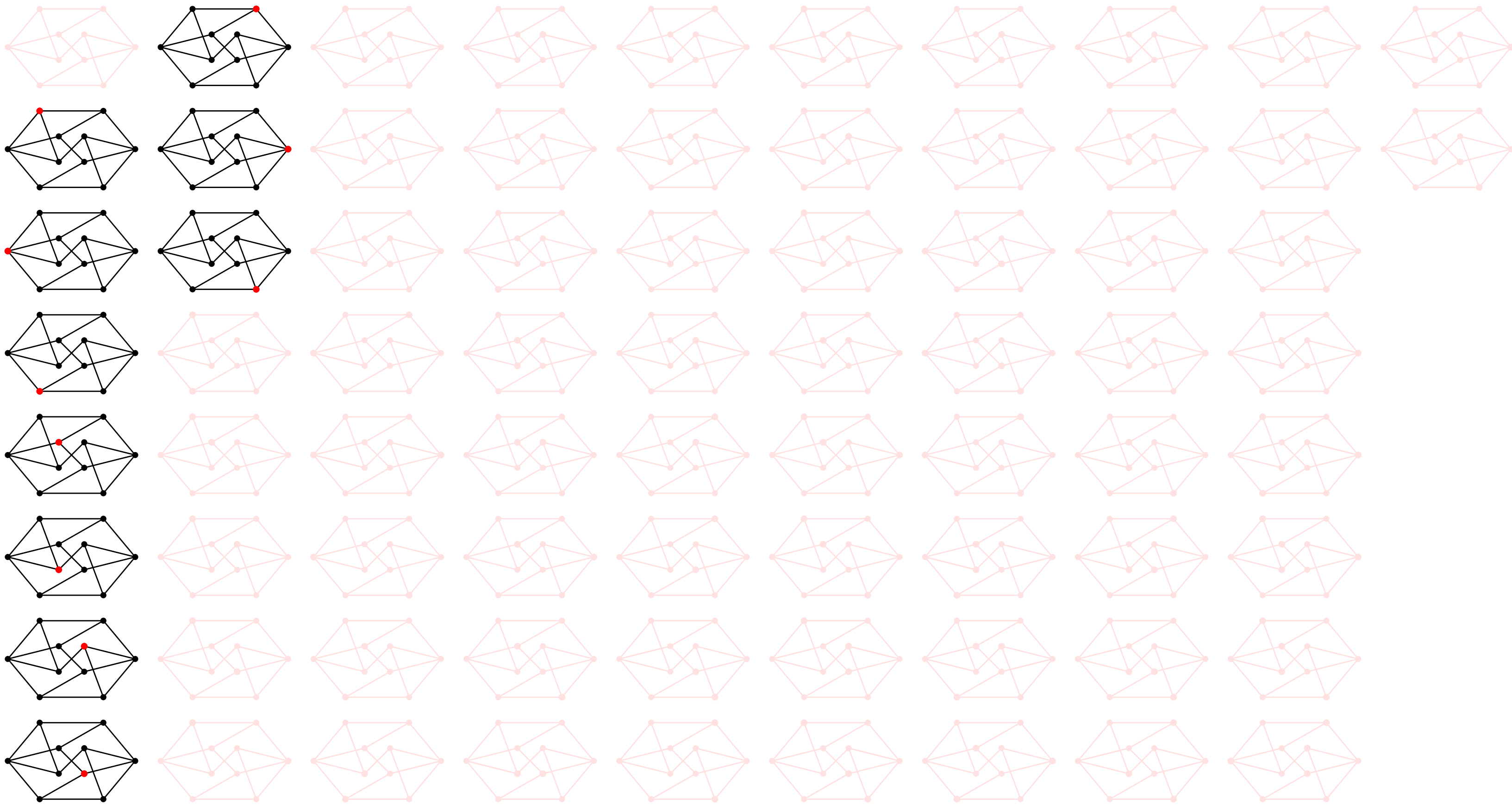
(5) # minimum maximal independent sets of G **1**

Independent sets of fixed size

74 independent sets

1 independent set of size 0

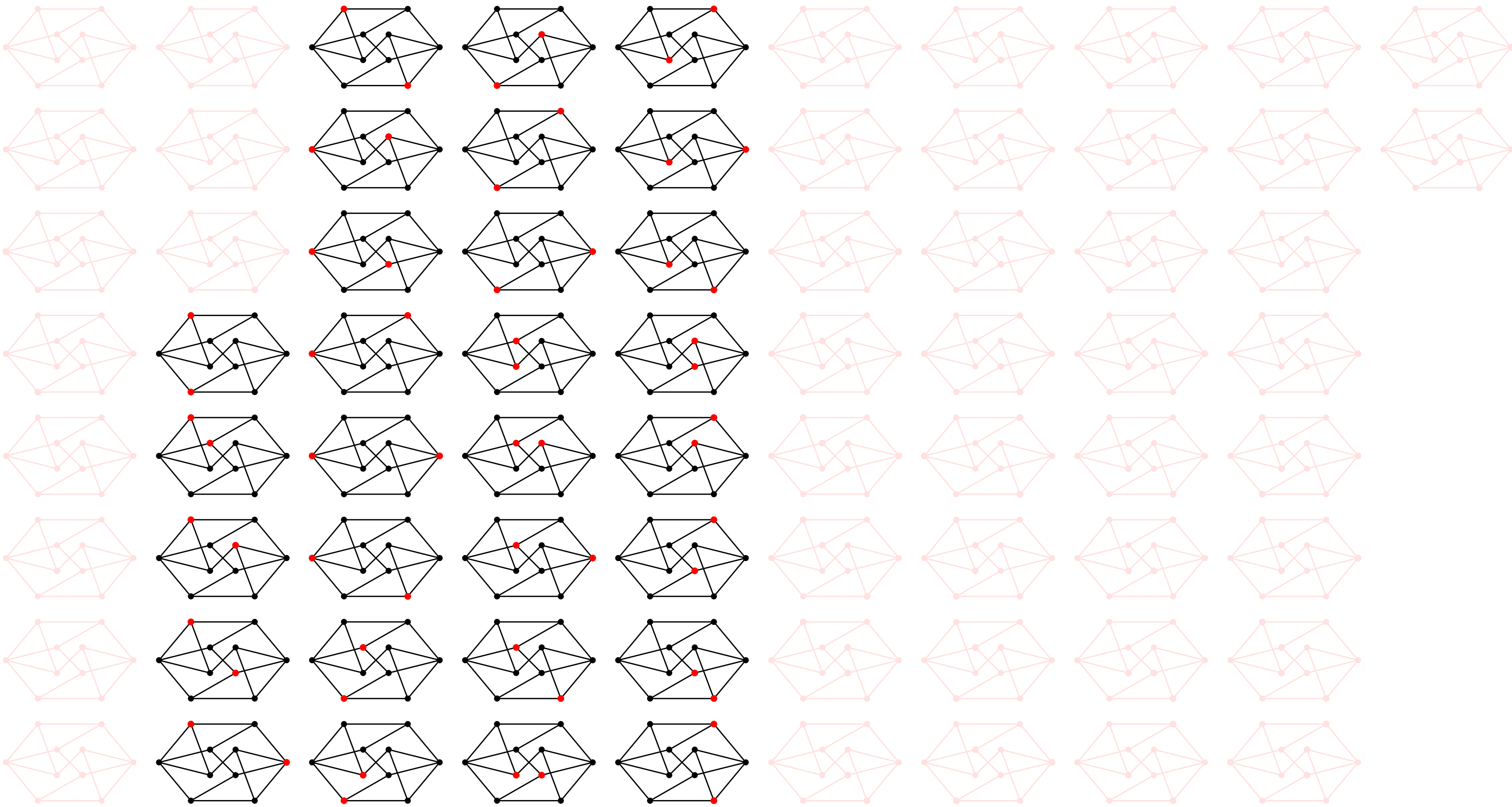
Independent sets of fixed size



74 independent sets

10 independent sets of size 1

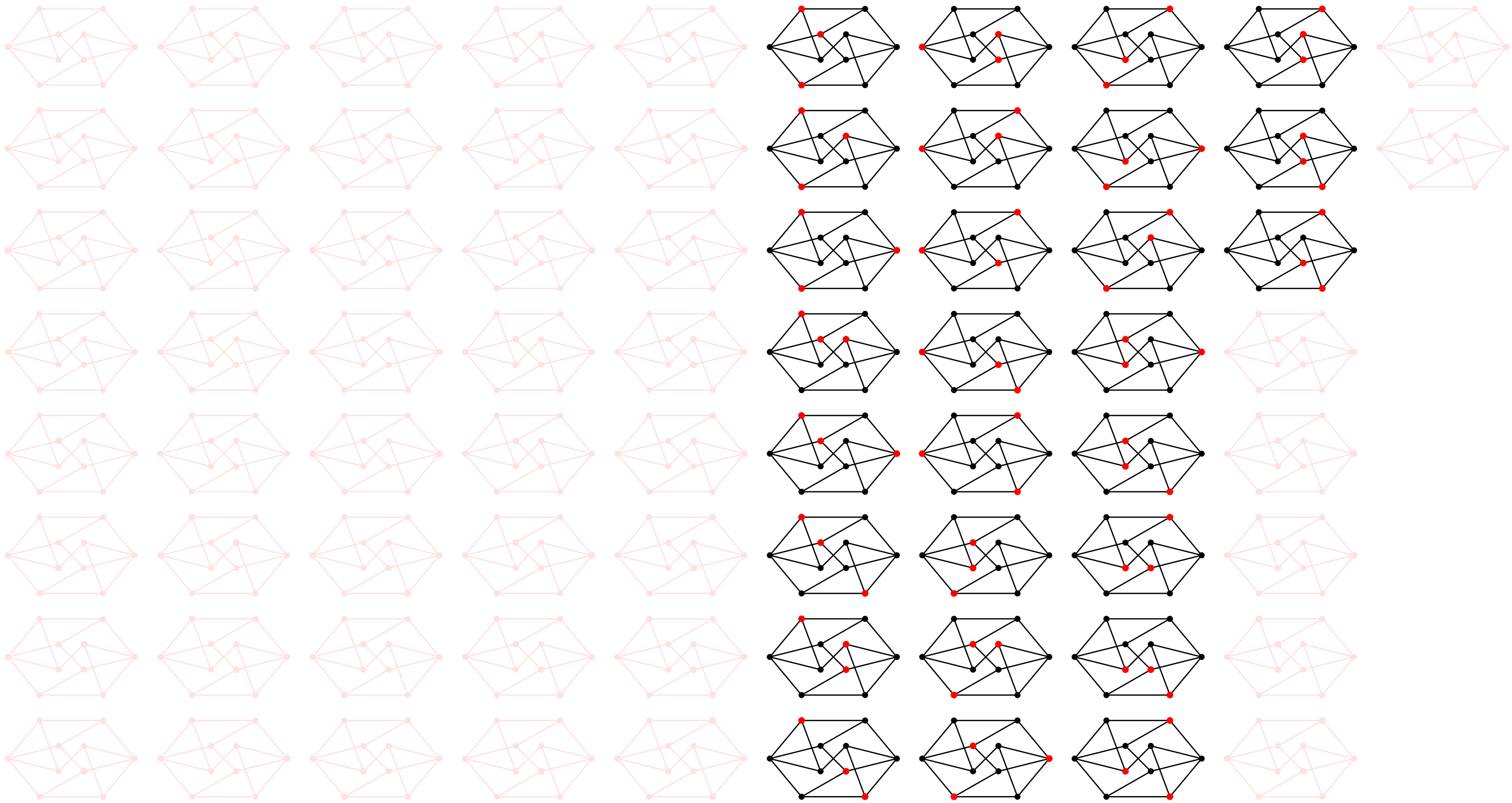
Independent sets of fixed size



74 independent sets

29 independent sets of size 2

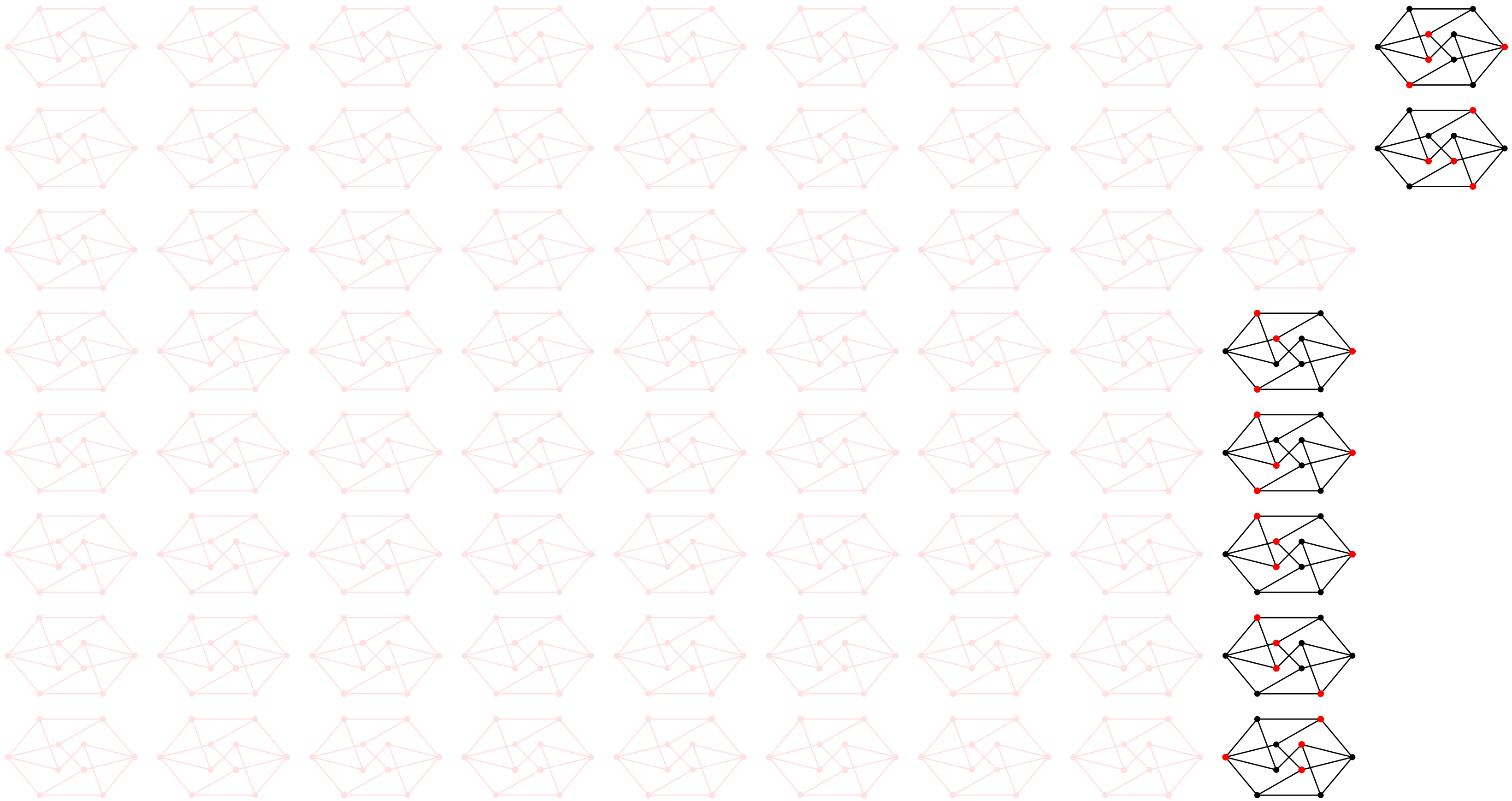
Independent sets of fixed size



74 independent sets

27 independent sets of size 3

Independent sets of fixed size



74 independent sets

7 independent sets of size 4

$\boxed{\text{Fact}}$ These counting problems are #P-complete (analoguous to NP-completeness).

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- the line graphs of bipartite graphs (Valiant '79)
- bipartite graphs (Provan & Ball '83)
- planar bipartite graphs of max deg 4 (Vadhan '01).

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- the line graphs of bipartite graphs (Valiant '79)
- bipartite graphs (Provan & Ball '83)
- planar bipartite graphs of max deg 4 (Vadhan '01).

$\Longrightarrow$ Focus on another class of perfect graphs!!

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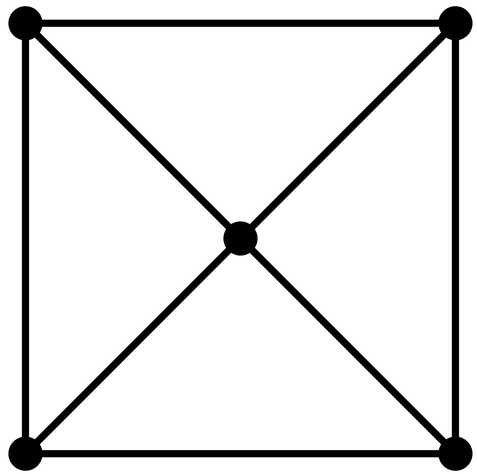

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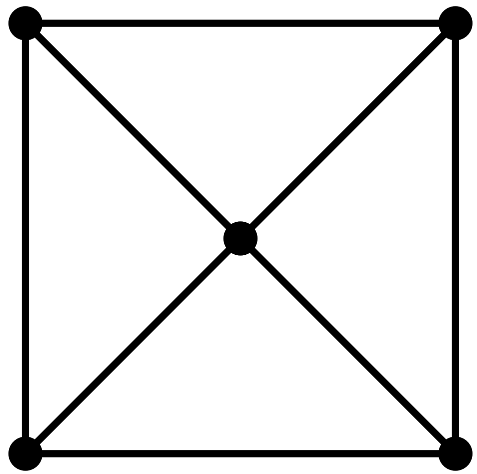

Not chordal

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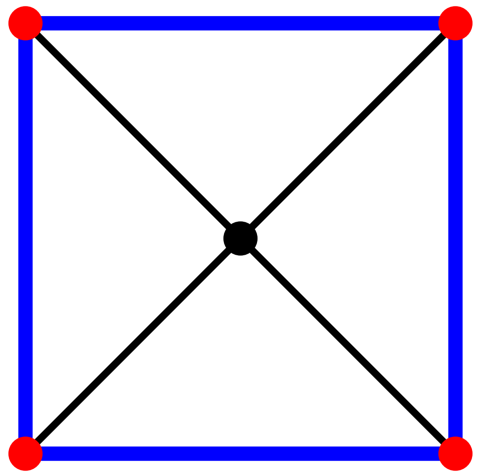

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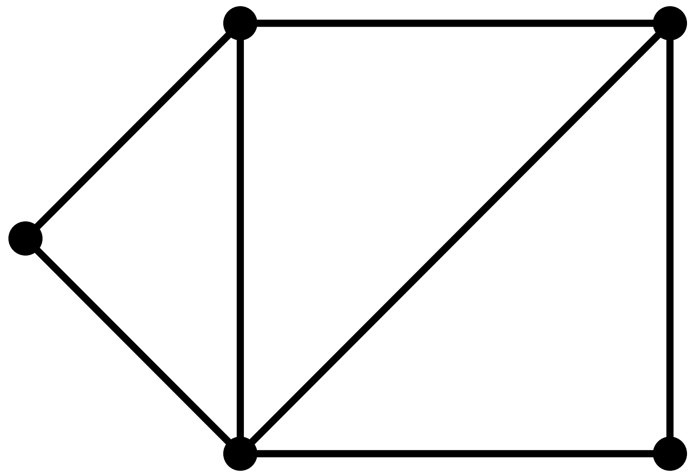

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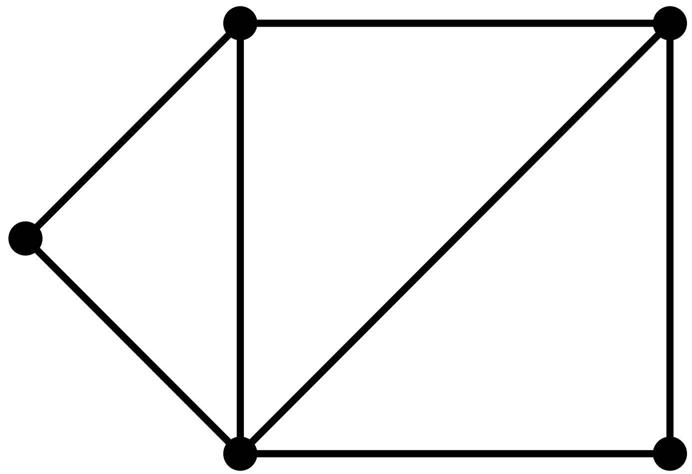

Chordal

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Output (1) # independent sets of G

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Rem "Finding one" is easy. (Gavril '72, Farber '82)

We obtain the following results for chordal graphs.

(1) # independent sets of G
(2) # maximum independent sets of G
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$O(|V| + |E|)$ alg.

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#P-complete

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#P-complete

Given $G = (V, E)$ a (connected) chordal graph

Def A **clique tree** of G is a tree T st

(1) the nodes of $T =$ the maximal cliques of G,

(2) $\forall\, v \in V$,

the nodes of T containing v induce a tree.

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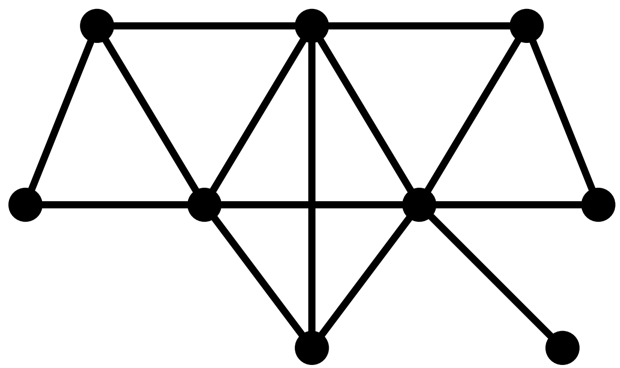

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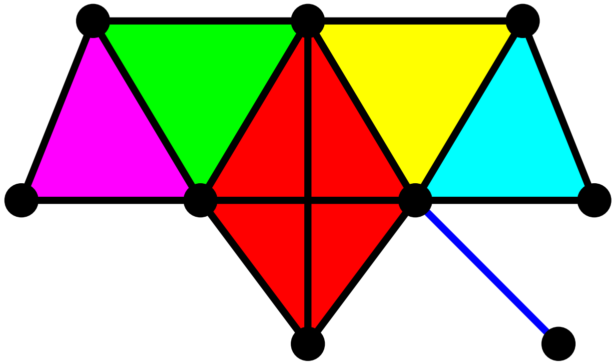

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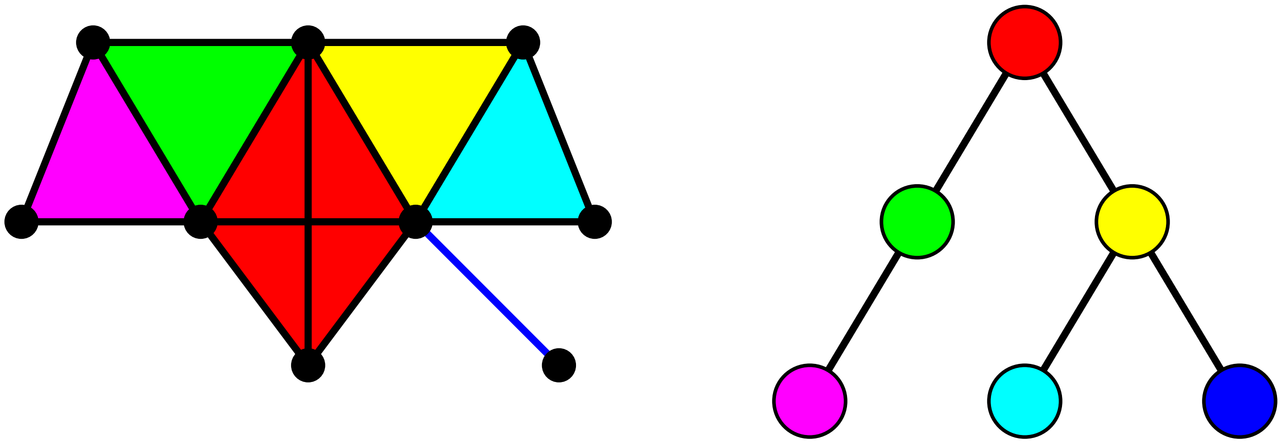

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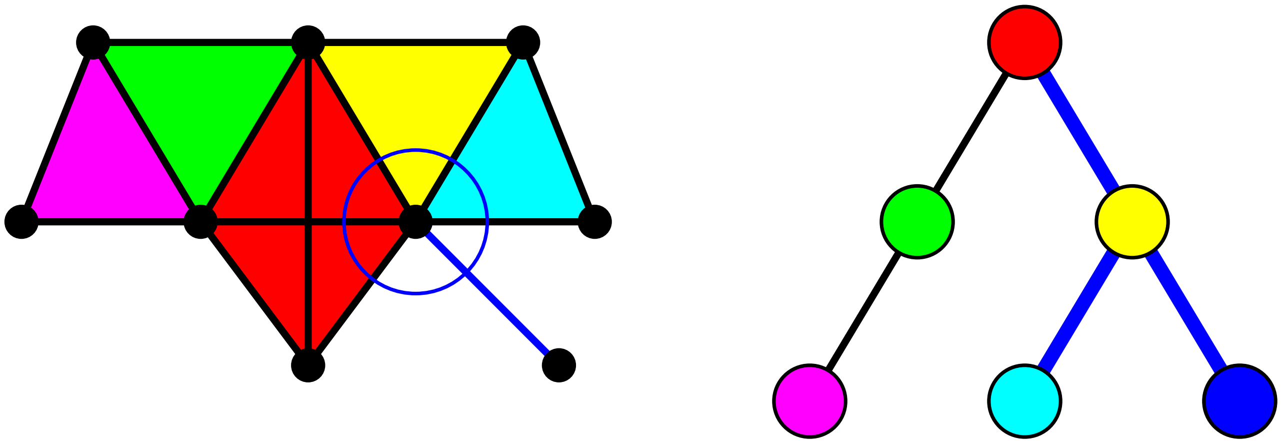

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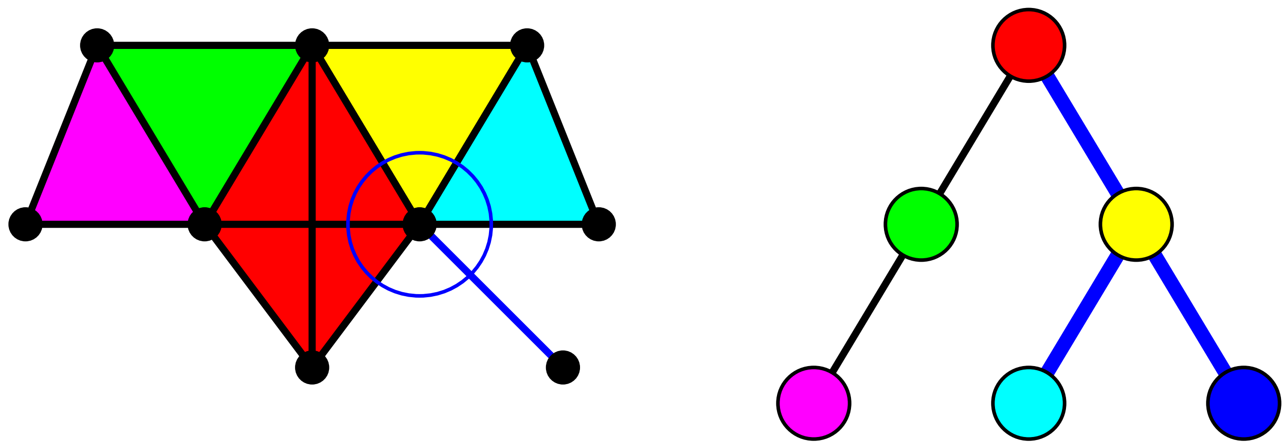

A clique tree can be computed in linear time.

Given $\quad G = (V, E)$ a (connected) chordal graph,

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T a clique tree of G, with root K

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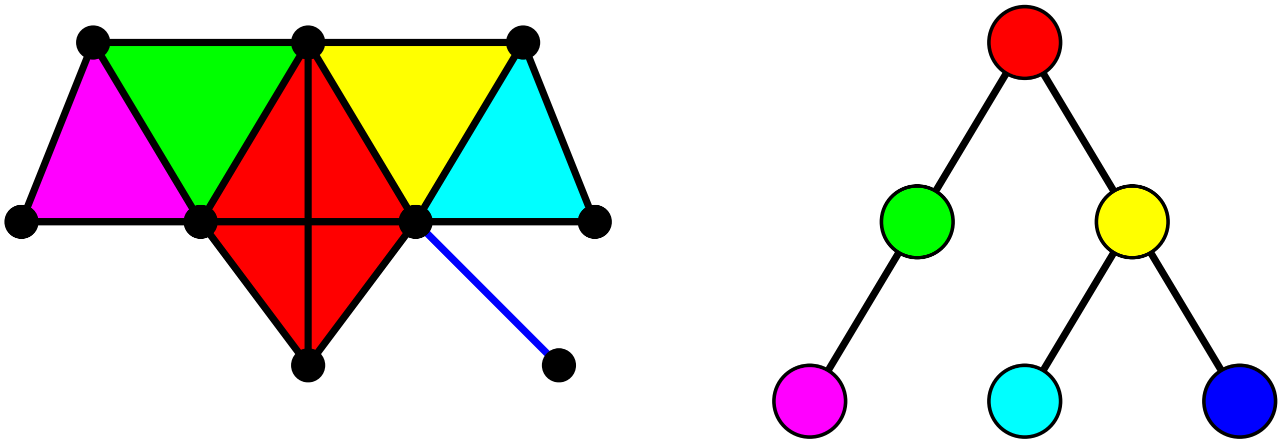
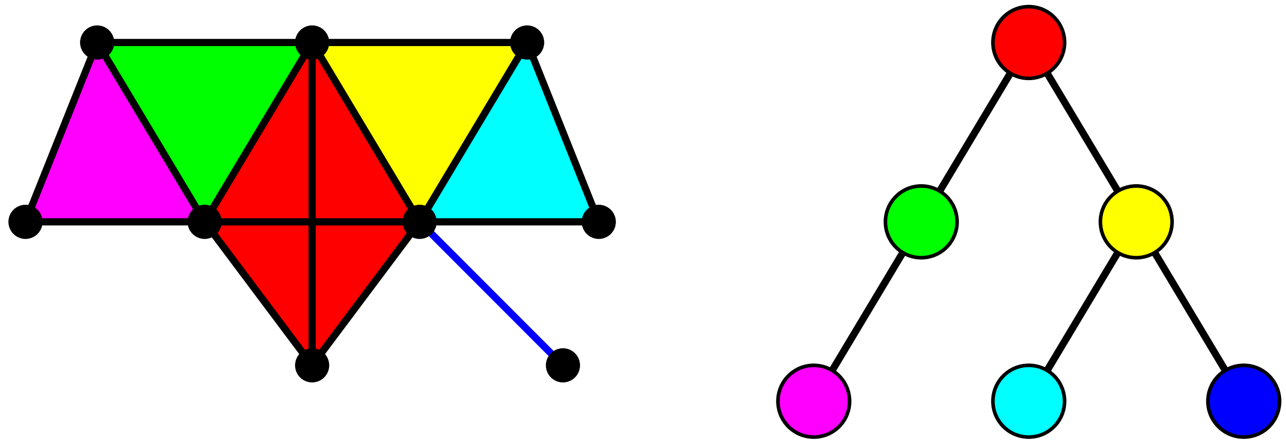


Given	$G = (V, E)$ a (connected) chordal graph, T a clique tree of G, with root K
Property	Every independent set of G either contains exactly one vertex in K or not.



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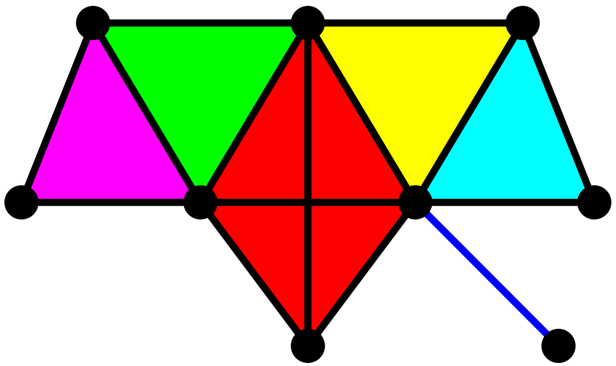
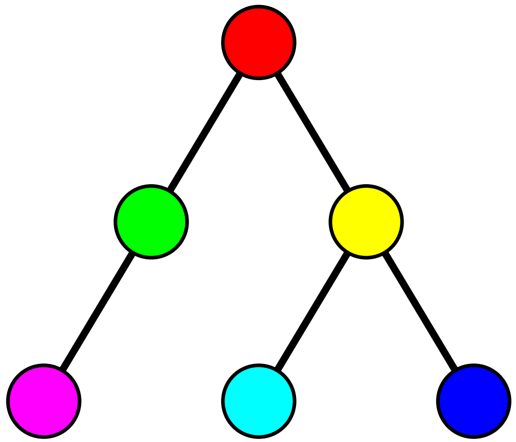
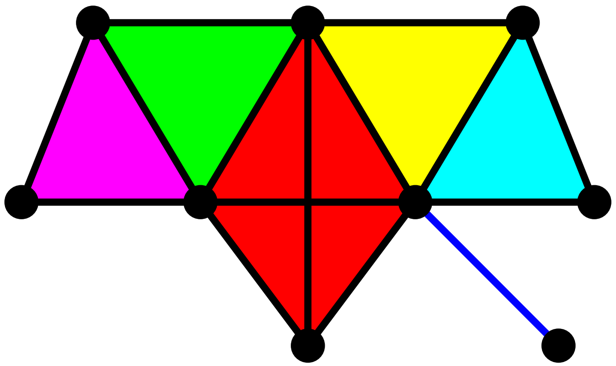


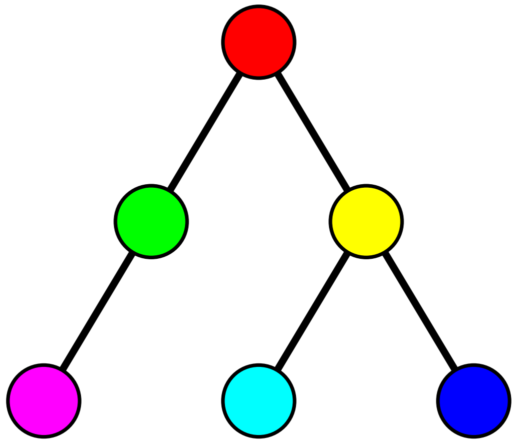


Leads to a recursive formula...

Given $G = (V, E)$ a (connected) chordal graph, T a clique tree of G, with root K

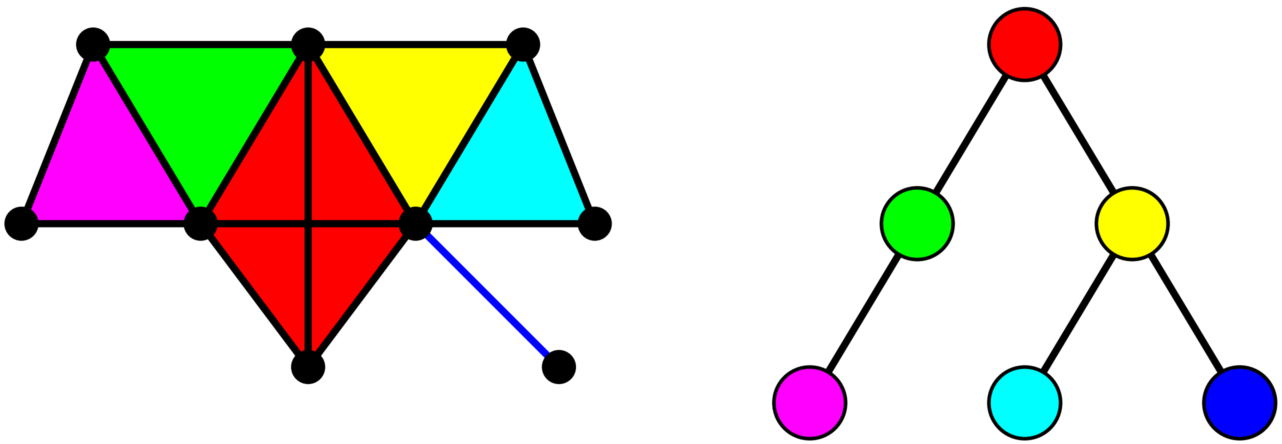
Lem # of independent sets of G =
# of independent sets of G excluding K +
$$\sum_{v \in K} (\text{\# of indepedent sets of } G \text{ containing } v)$$





Given $G = (V, E)$ a (connected) chordal graph,
T a clique tree of G with root K
K' a node of T

Notation $T(K')$ the subtree of G rooted at K'



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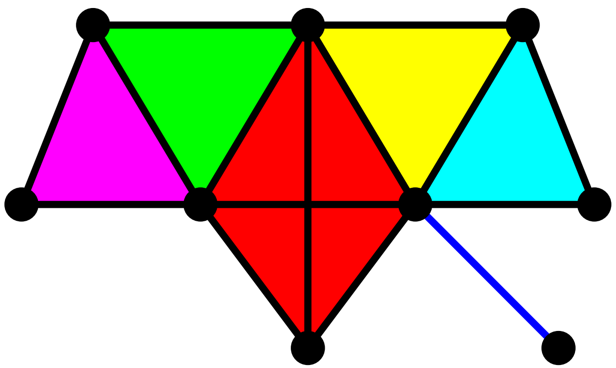

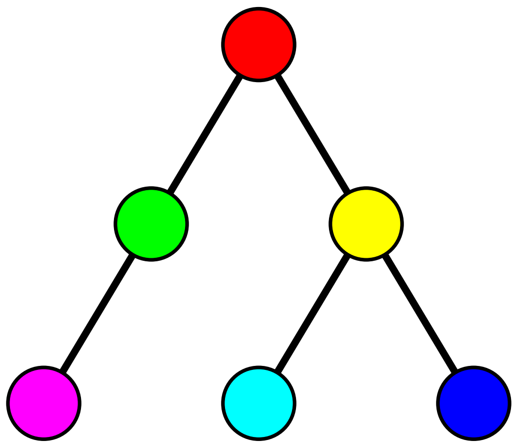

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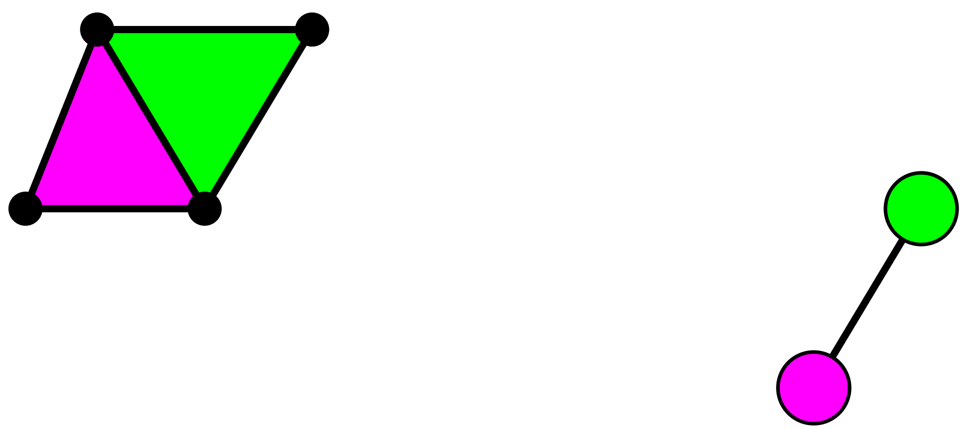

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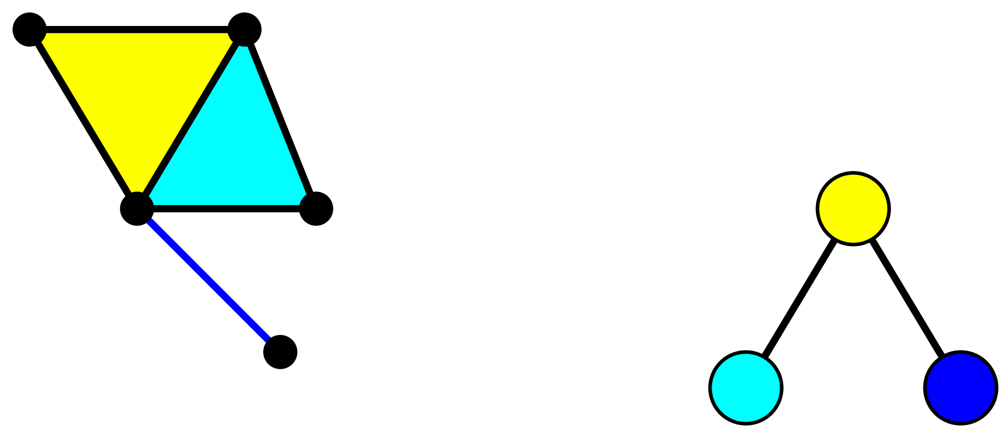
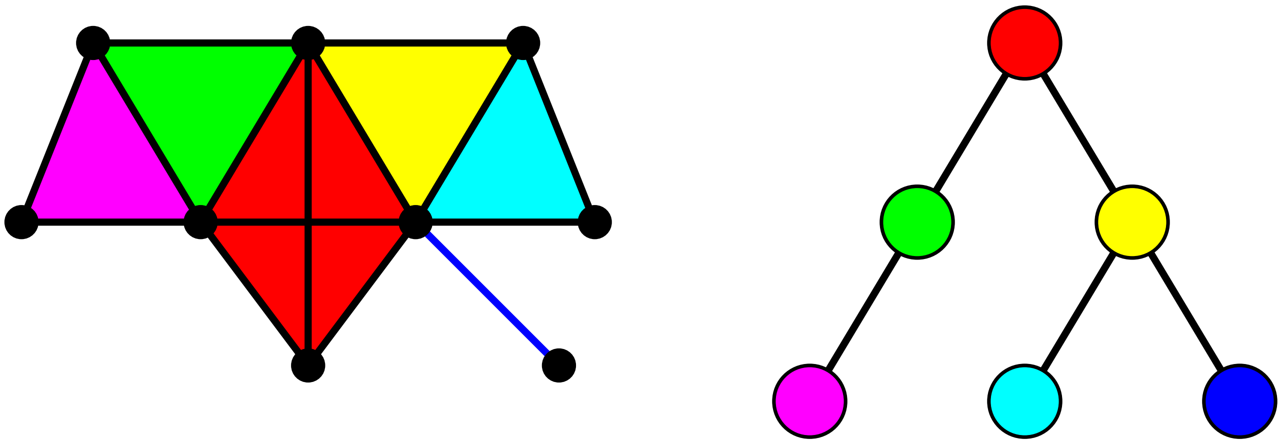


Given $G = (V, E)$ a (connected) chordal graph,
T a clique tree of G, with root K,
$K_1, \ldots, K_r$ the children of K in T

Lem # of independent sets of G excluding K =
$$\prod_i (\text{\# of independent sets of } G \text{ excluding } K \cap K_i)$$



Let G be a chordal graph and T be a rooted clique tree of G. For a maximal clique K of G which is not a leaf of the clique tree, let $K_1, \ldots, K_\ell$ be the children of K in T. Furthermore, let $v \in K$. Then, the following identities hold.

$$\mathcal{IS}(G(K)) = \overline{\mathcal{IS}}(G(K), K) \,\dot\cup \dot{\bigcup_{v \in K}} \mathcal{IS}(G(K), v);$$

$$\mathcal{IS}(G(K), v) = \{S \cup \{v\} \mid S = \bigcup_{i=1}^{\ell} S_i, S_i \in \left\{ \begin{array}{ll} \mathcal{IS}(G(K_i), v) & \text{if } v \in K_i \\ \overline{\mathcal{IS}}(G(K_i), K \cap K_i) & \text{otherwise} \end{array} \right\}\};$$

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where $\mathcal{IS}(G)$ denotes the family of independent sets in G, $\mathcal{IS}(G, v)$ denotes for a vertex v the family of independent sets in G including v, $\overline{\mathcal{IS}}(G, U)$ denotes the family of independent sets in G including no vertex of U for a vertex set U.

Recursive formula

Let G be a chordal graph and T be a rooted clique tree of G. For a maximal clique K of G which is not a leaf of the clique tree, let $K_1, \dots, K_\ell$ be the children of K in T. Furthermore, let $v \in K$. Then, the following identities hold.

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A detailed analysis yields a linear-time algorithm to count the independent sets in a chordal graph!

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Thm The independent sets in a chordal graph $G = (V, E)$ can be counted in $O(|V| + |E|)$ time.

We obtain the following results for chordal graphs.

(1) # independent sets of G
(2) # maximum independent sets of G
(3) # independent sets of G of fixed size

$O(|V| + |E|)$ alg.

(4) # maximal independent sets of G
(5) # minimum maximal independent sets of G

#P-complete

We use the following counting problem
(a counting version of the set cover problem).

Given	X a finite set, $\mathcal{S} \subseteq 2^{X}$ a family
Output	The number of subfamilies $\mathcal{S}' \subseteq \mathcal{S}$ st $\bigcup_{Y \in \mathcal{S}'} Y = X$.

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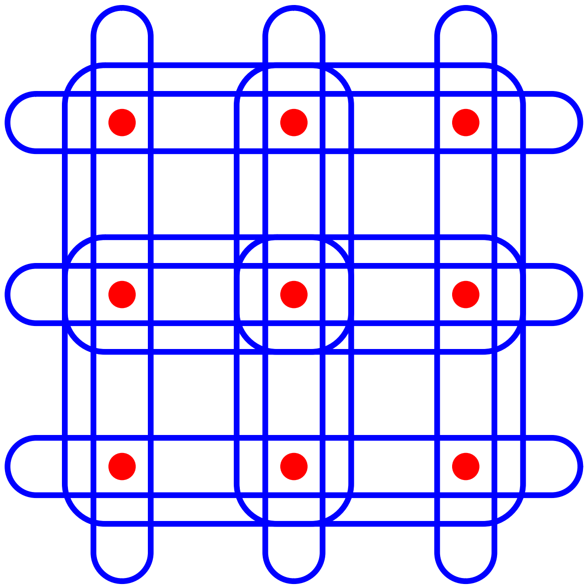

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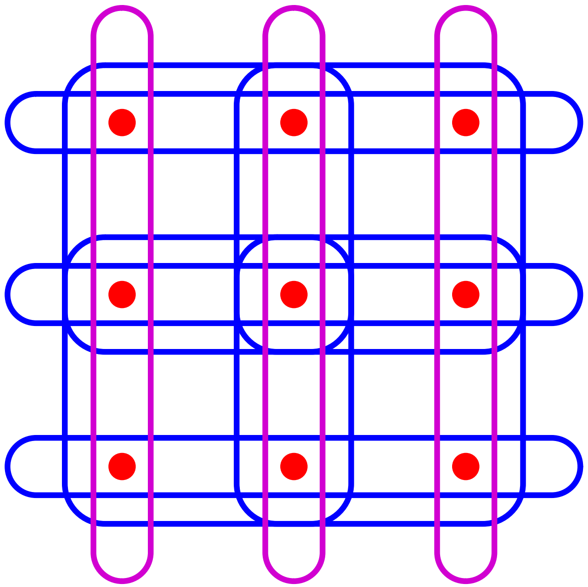

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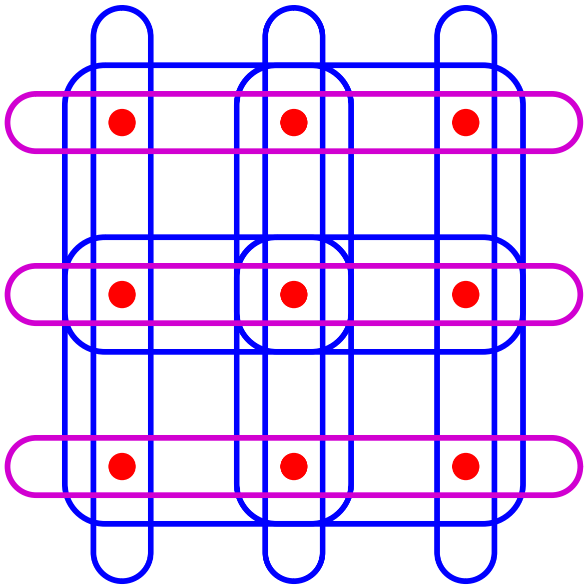

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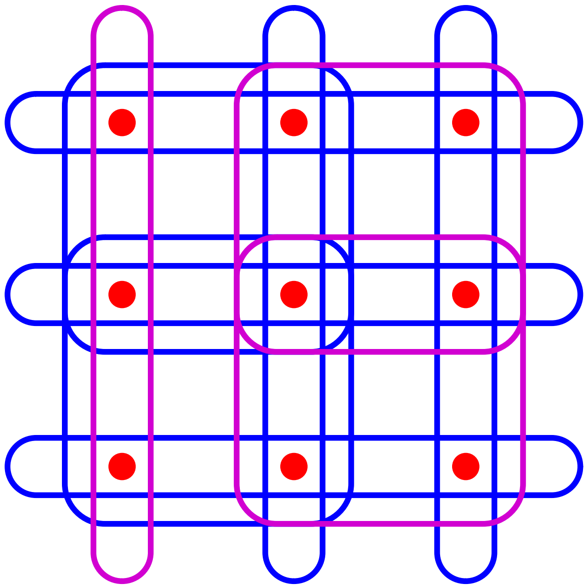

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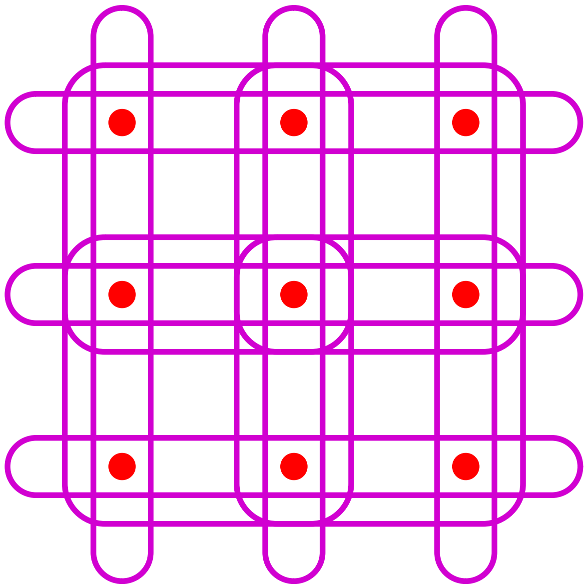

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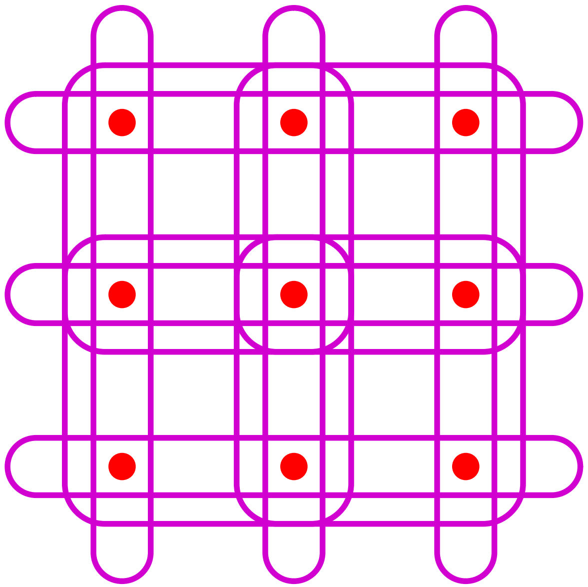
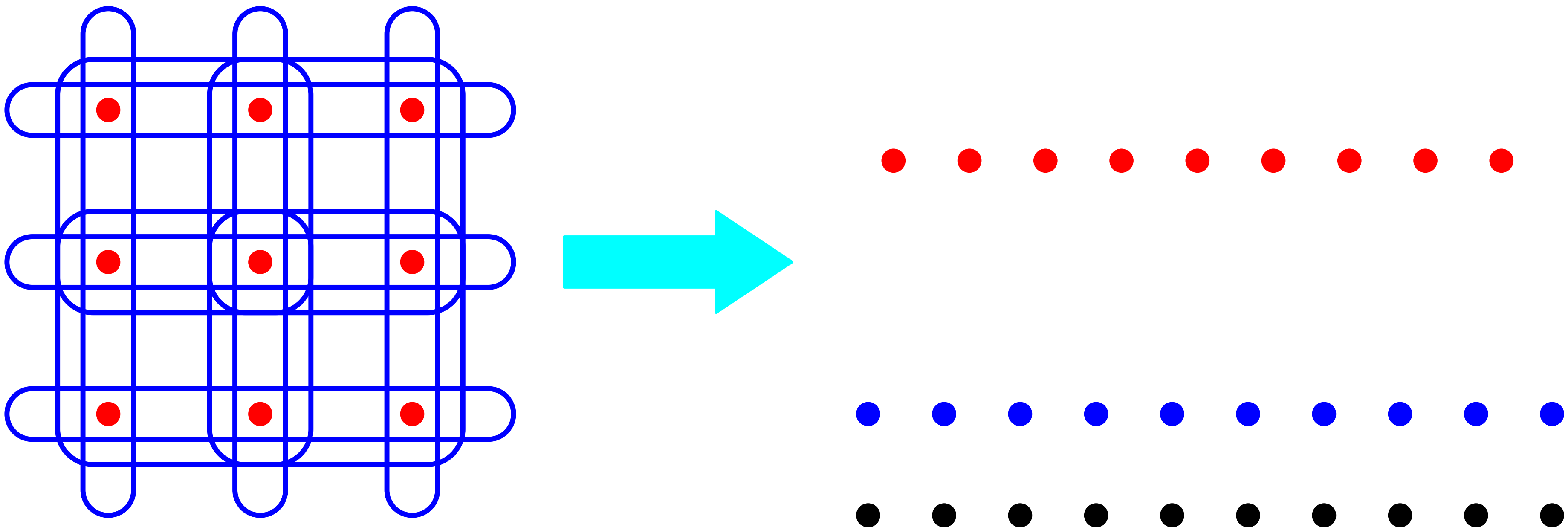
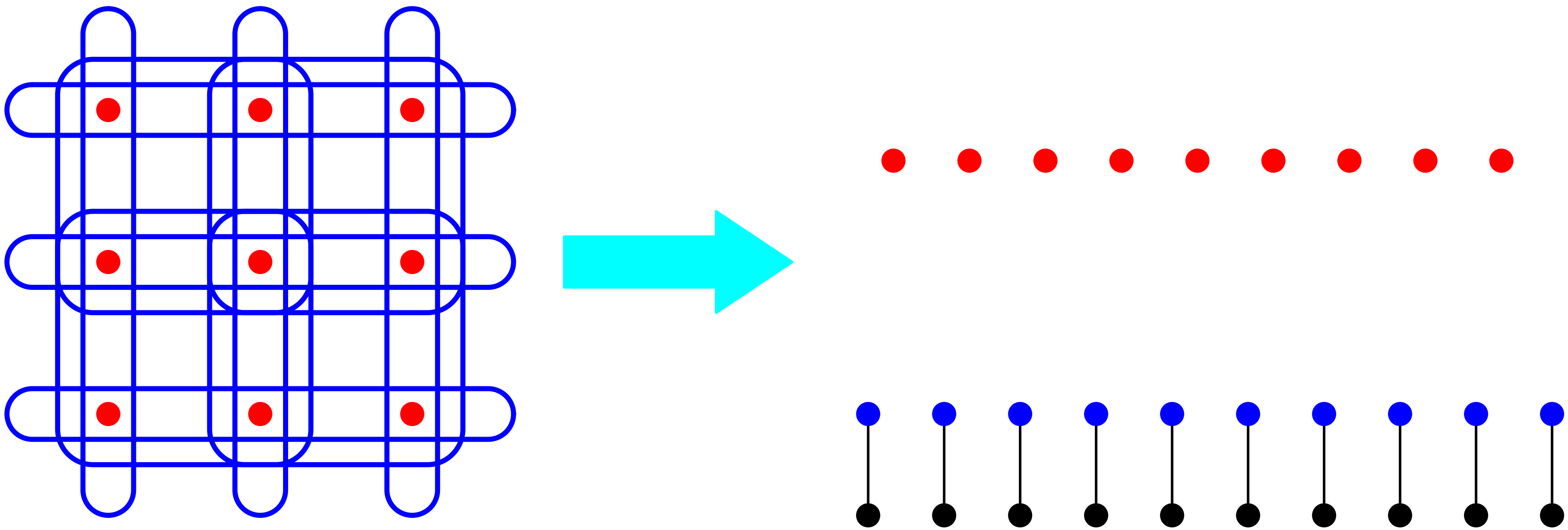


Actual reduction





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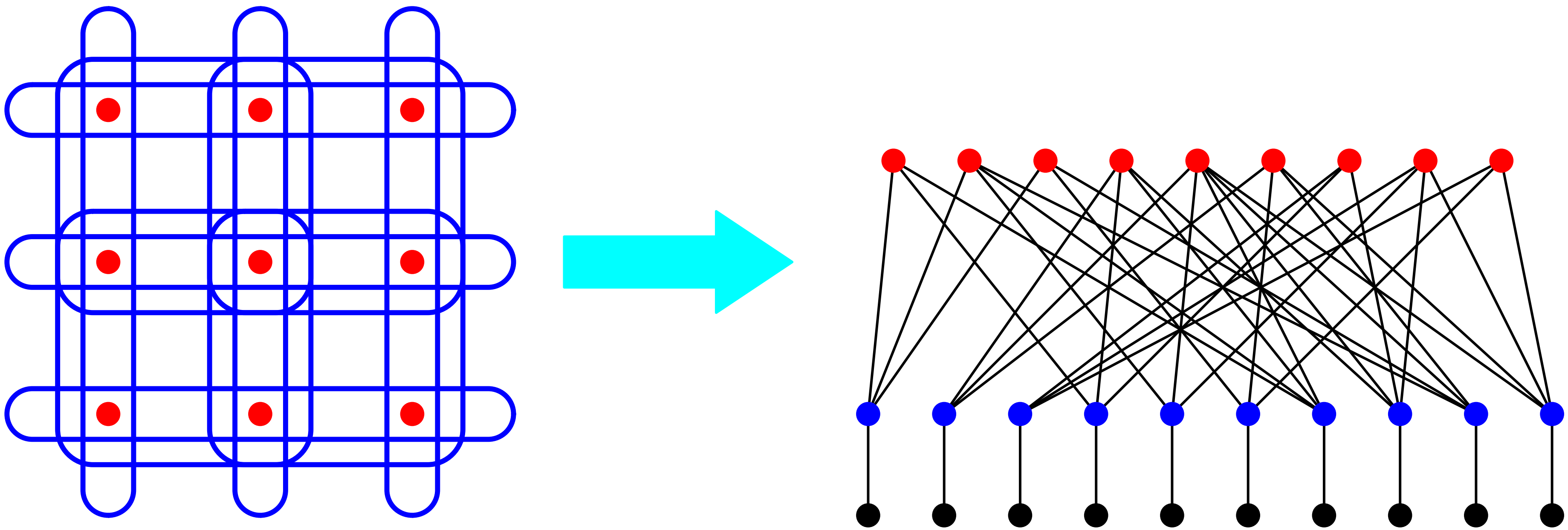

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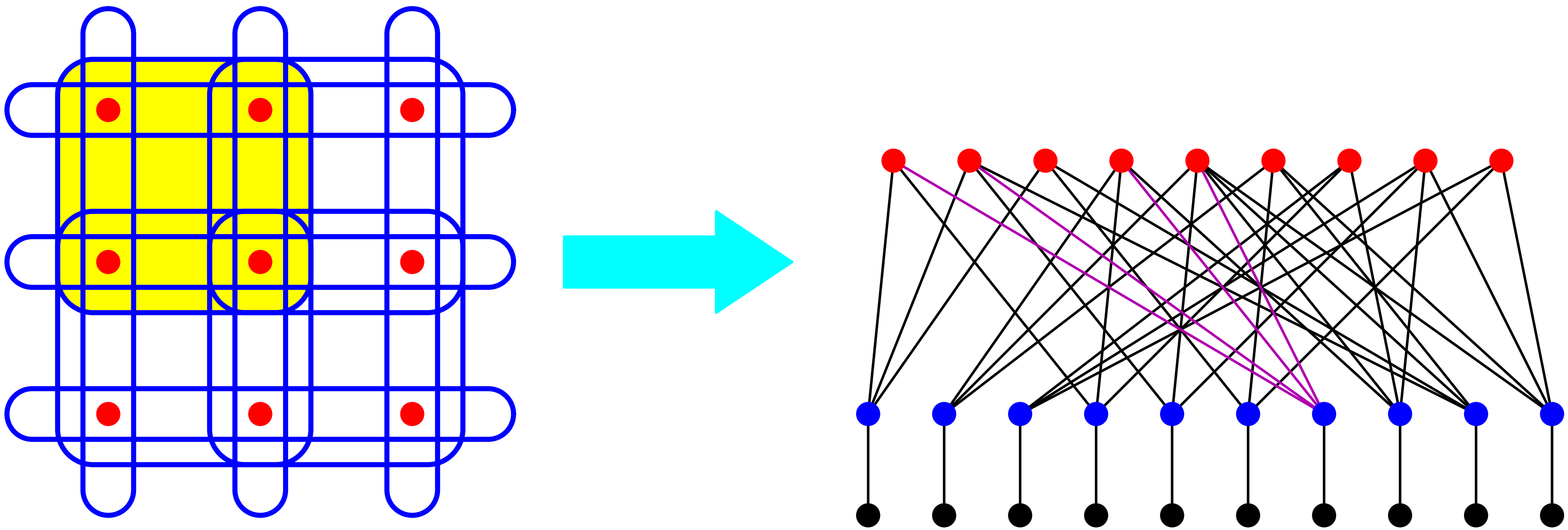

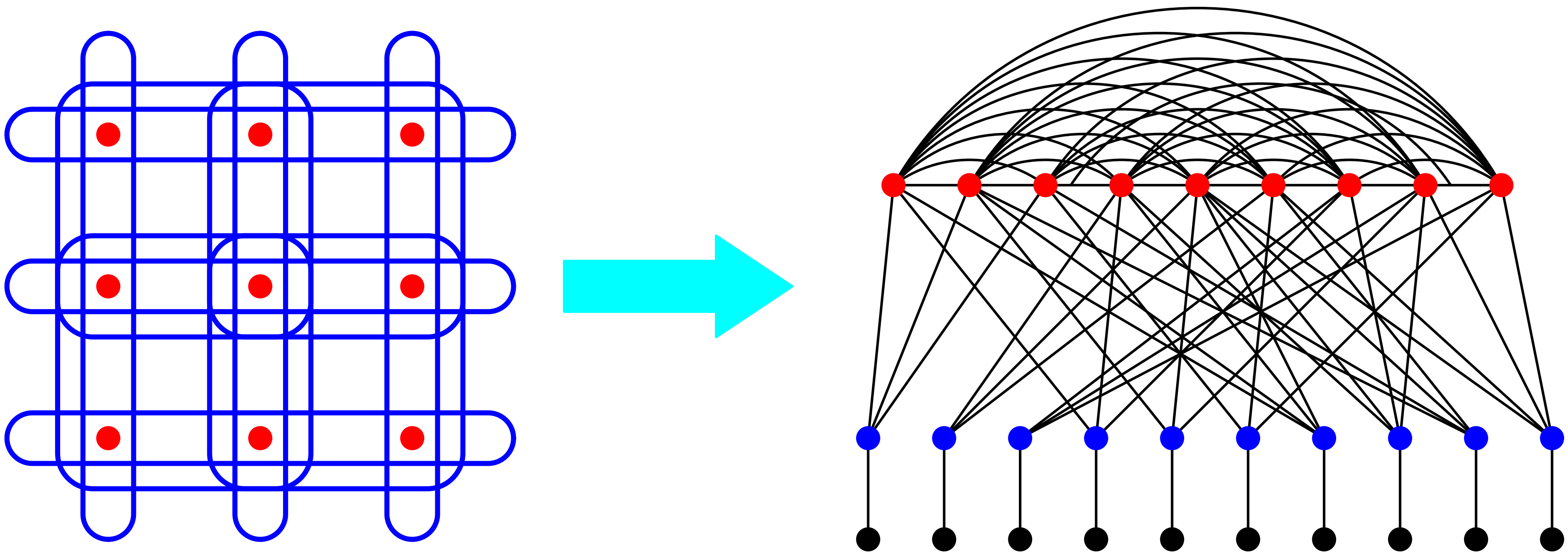

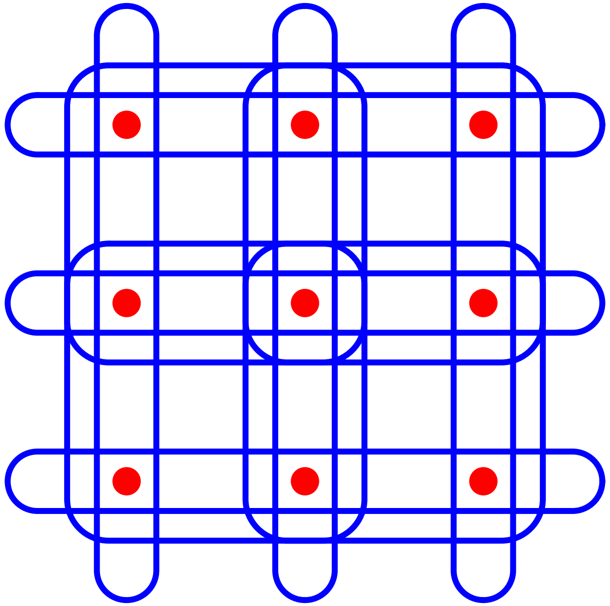

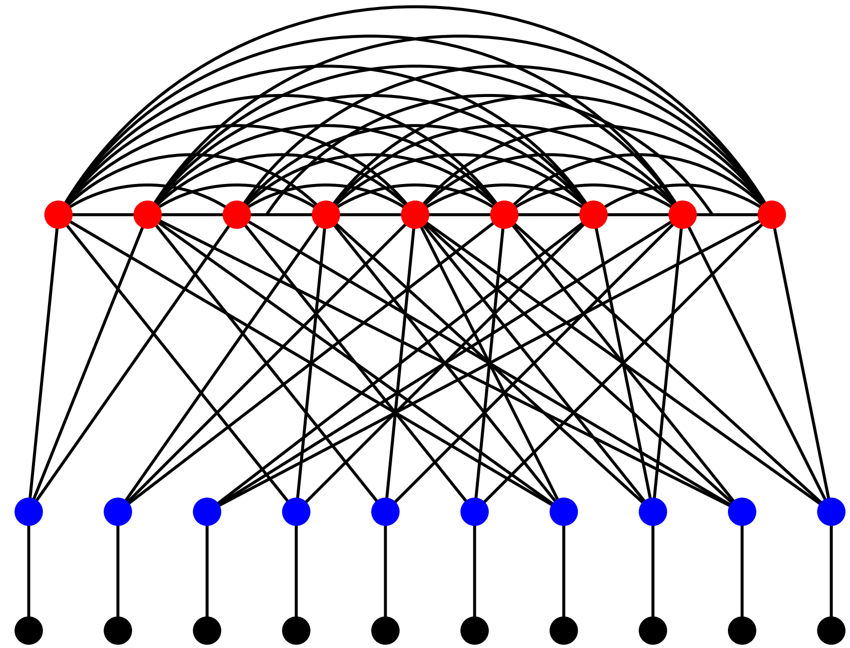

$$\# \text{ set covers} + \sum_{x \in X} 2^{\#\text{ sets not containing } x}$$

$$\|$$

$$\# \text{ maximal independent sets}$$

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[End of my talk]

Merci beaucoup.